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### Potential Vulnerabilities of Timbuktu Remote Control Software

This document is intended to explain the functionality and potential vulnerabilities of the Windows Remote Access/Remote Control Software application called Timbuktu. It has been said that "A picture is worth a thousand words." This is certainly the case when working with end-users who experience problems using a particular piece of software or perform some process with their Windows workstation. Using remote control software to identify the specific error messages or observe a series of dialog boxes can be dramaticall...

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# Potential Vulnerabilities of Timbuktu Remote Control Software

David Batz

October 9, 2002

## Abstract:

The Problem: In today's connected world, there is often a need for help desk personnel or support staff to access end-user Windows workstations or remote servers. Many times, the devices being accessed may be geographically dispersed (hundreds or thousands of miles away) from support personnel.

Timbuktu – A Solution to the Problem? Disclaimer: The reader should understand that although this paper is neither for nor against the use of Timbuktu software as a Windows Remote Access /Remote Control solution, there are a number of potentially serious vulnerabilities that may be encountered through the use of the product. After reviewing the issues and potential responses the reader will be better equipped to make an informed decision about remote control products in general, and Timbuktu specifically.

## Introduction:

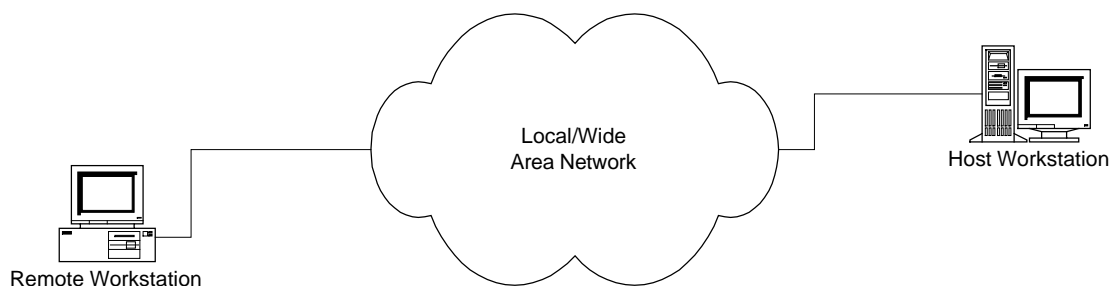
This document is intended to explain the functionality and potential vulnerabilities of the Windows Remote Access/Remote Control Software application called Timbuktu.

It has been said that "A picture is worth a thousand words." This is certainly the case when working with end-users who experience problems using a particular piece of software or perform some process with their Windows workstation. Using remote control software to identify the specific error messages or observe a series of dialog boxes can be dramatically more efficient than attempting to wade through fuzzy or inaccurate descriptions of problems.

There are a number of technical solutions available to address the challenge of remote access to windows hosts. These solutions include: Symantec PC Anywhere, VNC, Citrix, and Windows Terminal Service. This paper is concerned with Timbuktu Pro 2000 (Version 2.0 Build 815 Es) deployed on Windows 2000 devices (Professional and Server.) The product is developed by Netopia Inc., <http://www.netopia.com/en-us/software/products/tb2/index.html>.

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## Communications Overview



Client software on the remote workstation connects to a set of services on the host. In order to provide for total unattended connectivity on the Timbuktu host, several programs are started as services. A detailed explanation of the connection process will be covered later in the paper.

Remote access privileges to the host machine are controlled by a single file; tb2.plu. By default, this file is located in the following directory structure:  
C:\Program Files\Timbuktu Pro. -1 (Blue Boar-insecure.org)

## Communication Services

Timbuktu software supports seven flavors of functionality:

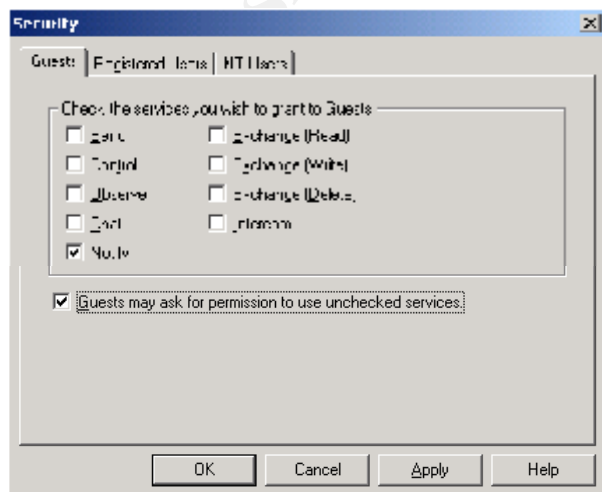
- **Send** allows users to send messages with attached files and folders to a remote host
- **Exchange** allows users to access the hard drive of a remote Timbuktu Pro host — even a Macintosh. Authorized users can copy and move files and folders between computers and delete files from either computer
- **Control** allows users to control a remote host from their own desktop
- **Observe** allows users to observe a remote host without controlling it
- **Notify** alerts users when a remote Timbuktu Pro host becomes active.
- **Chat** allows users to carry on a keyboard conversation with a remote user
- **Intercom** allows users to speak directly to a remote user through a host's audio hardware

Timbuktu supports three different types of users

- Guest Users (which include Ask for Permission Users and Temporary Guests)
- Registered Users
- Windows Domain Users. -2 (Netopia Help File)

On a Windows 2000 machine, a graphical user interface is used to configure Timbuktu User-Ids, passwords, and user access rights.

## Guest Access



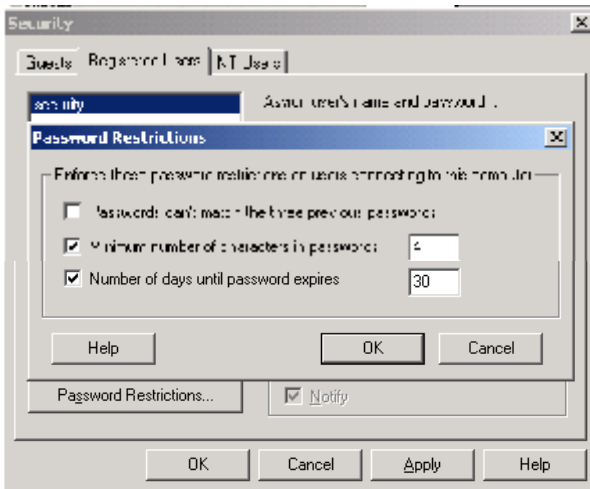
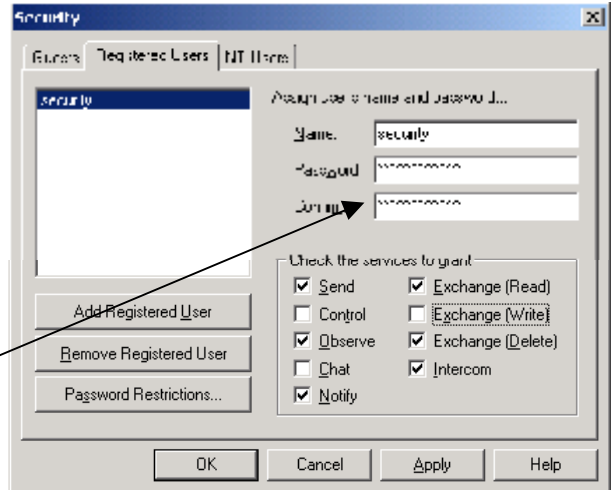
In Timbuktu, Guests function as a global “Everyone” group. Any rights granted to Guest users are implicitly granted to all workstations with the Timbuktu software. Rights granted to the “Guests” group may **not** be removed from “Registered Users” or “NT Users.” **Extreme caution should be used when defining “Guest” access.** Guest privilege information is stored in the following directory structure: C:\Program Files\Timbuktu Pro\tb2.plu.

## Registered User Access

In this dialogue, the system administrator can define specific users, and their associated privileges. Registered User names, encoded passwords, and privileges are also stored in the tb2.plu file.

This file will be examined in detail later in this paper.

Although the password entry field looks quite large, the maximum password length is 15 characters.



Timbuktu does not support a “Hard” lockout in the event that a Registered User supplies an incorrect password repeatedly. An attacker can keep banging on a host machine.



The system administrator can also define *some* Password Restrictions for Registered Users.

As can be seen in the dialog box, the abilities to enforce specific password restrictions (policies) are limited to:

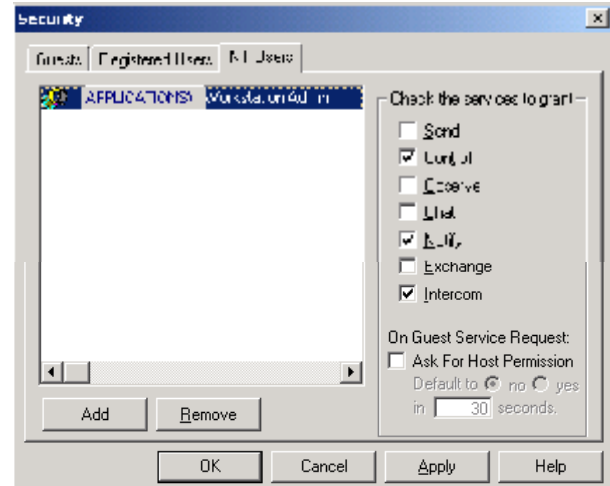
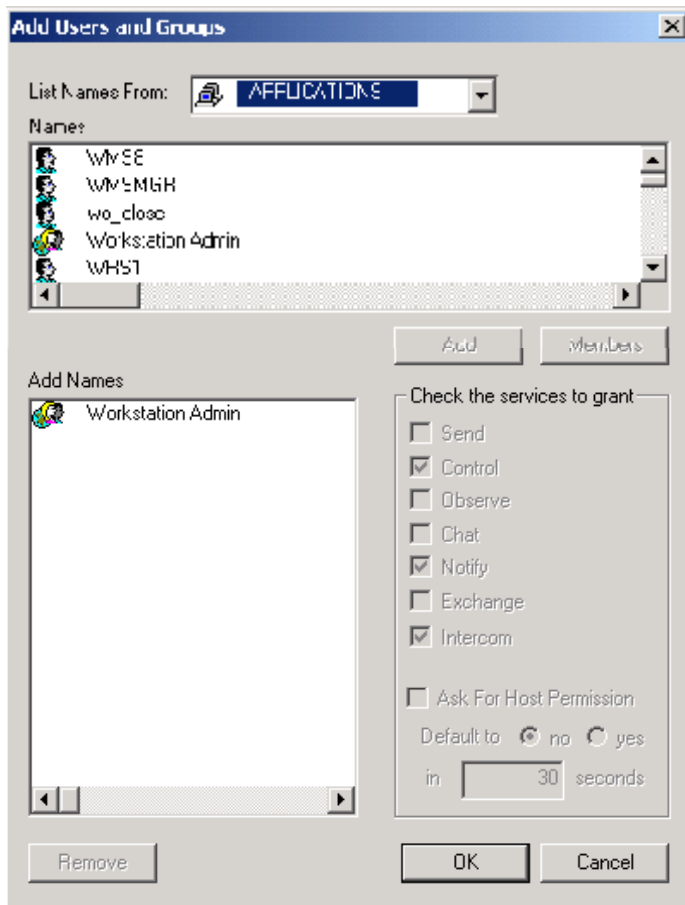
- A password may not match any of the previous three passwords
- Passwords must be a minimum length (the application does not support password lengths greater than 15 characters)
- Passwords can be set to expire in the future

Timbuktu does **not** prevent “weak” passwords:

- Dictionary words (including foreign and technical dictionaries)
- Anyone's or anything's name
- A person, place or thing
- A proper noun
- A phone number
- Passwords of the same character (e.g. aaa)
- Simple pattern of letters on keyboards (e.g. qwerty)
- Any of the above reversed or concatenated (e.g. ytrewq)
- Any or the above with digits prepended or appended (e.g. aaa1)

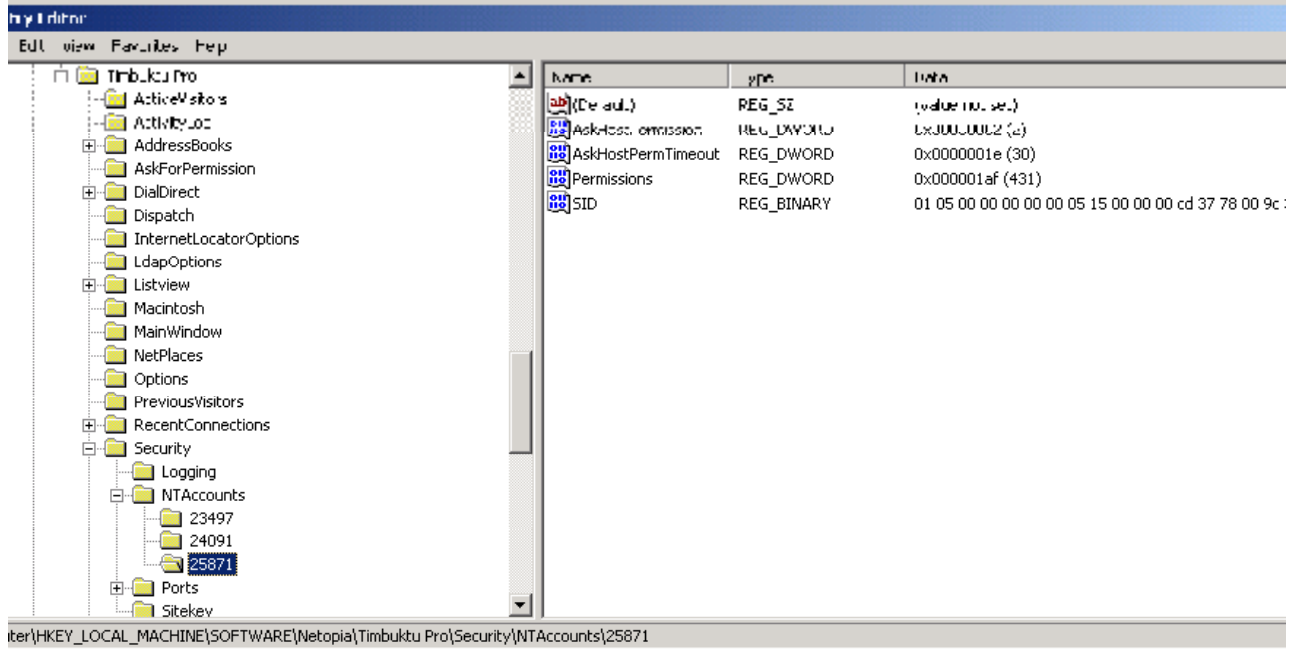
-3 (Maui High Performance Computing Center Kerberos Password Policy)

## NT User Access



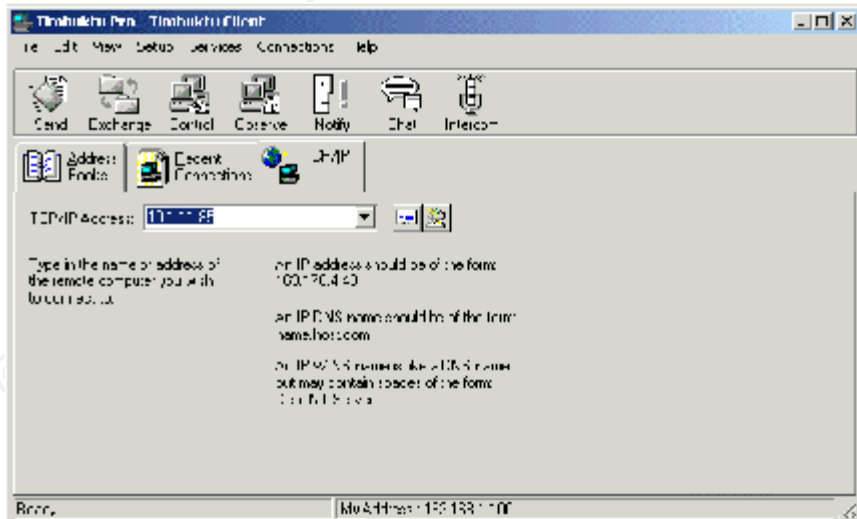
If the Timbuktu host is part of a NT Domain, the system administrator may add individual NT user or group account access. In this situation, there is no opportunity to use a password policy specific to Timbuktu. The password policy of the NT domains is inherited, and used by Timbuktu. When NT Users are selected, information about the specific User ID and rights attributes are placed into the NT registry, rather than a configuration file in the Timbuktu sub-directory.

This graphic shows the Registry entries for a Timbuktu configuration with three NT Users identified. Each User has their own registry Key defined under the NTAccounts Key.



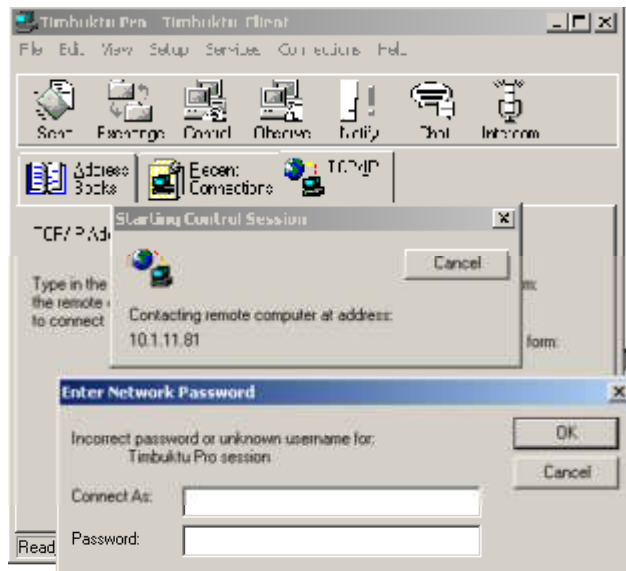
**Using the product:**

The Graphical User Interface of Timbuktu (as a client) has the following appearance:

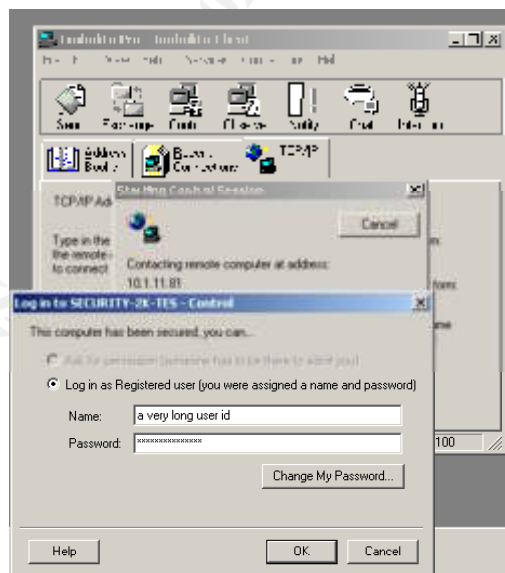


After typing a computer name or address, the end-user clicks on the appropriate function, such as Send, Exchange, Control, etc.

On Windows 2000 platforms, the Timbuktu Host software assumes that the primary application authentication will be based on Windows Domain membership. When this is not the situation, the client will be shown the following dialog:



“Registered and Guest Users” must press the Escape key to encounter the following dialog:



At this point, the client can enter the Registered User credentials, or if the host system is configured, the client can ask for permission to connect.

Once the Timbuktu client is connected to the host, the client inherits the rights and permissions of the end-user who is logged into the host system. If the Timbuktu host is logged in as an administrator, the remote client gains those privileges.

Let's take a look at what happens under the covers.

After a default installation, a Timbuktu host will be running the following programs:

Name	Memory Used
Tb2RCAssist.exe	1,312 K
Tb2pro.exe	13,796 K
Tb2logon.exe	1,532 K
Tb2launch.exe	884 K

After these programs have started, but before a remote session is established, a Windows Timbuktu host will listen on port 407/UDP.

After a remote session is established, a Windows Timbuktu host will listen on the following ports, depending on the services negotiated during session startup:

Service	Host Listening Port
Control	1417 /TCP
Observe	1418 /TCP
Send Files	1419 /TCP
Exchange Files	1420 /TCP
Chat	Dynamic TCP Ports
Notify	Dynamic TCP Ports
Intercom	Dynamic TCP and UDP Ports
Ask for Permission	Dynamic TCP and UDP Ports

-4 (Netopia FAQ Article)

The Following Windump output shows the communications between the Timbuktu Client and Host. This example shows an initiation of a **“Control”** session:

```
22:47:11.358305 timb-client.1540 > Timbuktu-host.407: udp
22:47:11.363389 Timbuktu-host.407 > timb-client.1540: udp
22:47:15.324970 timb-client.1541 > Timbuktu-host.407: udp
22:47:15.330155 Timbuktu-host.407 > timb-client.1541: udp
22:47:15.332263 timb-client.1542 > Timbuktu-host.407: udp
22:47:15.341369 Timbuktu-host.407 > timb-client.1542: udp
22:47:15.342409 timb-client.1543 > Timbuktu-host.1417: S 3257468345:3257468345(0)
win 16384 <mss 1460,nop,nop,sackOK> (DF) [Now, the initial Syn]
22:47:15.342833 Timbuktu-host.1417 > timb-client.1543: S 3281877292:3281877292(0)
ack 3257468346 win 17520 <mss 1460,nop,nop,sackOK> (DF) [Syn, Ack]
22:47:15.342870 timb-client.1543 > Timbuktu-host.1417: . ack 1 win 17520 (DF) [Ack –
The TCP communications channel is now open,]
```

**First two attempts to exchange credentials fail**

**Third try is the charm**



22:47:15.343240 timb-client.1543 > Timbuktu-host.1417: P 1:17(16) ack 1 win 17520 (DF)  
**[The Timbuktu session starts]**

22:47:15.475504 Timbuktu-host.1417 > timb-client.1543: . ack 17 win 17504 (DF)

22:47:15.482306 timb-client.1543 > Timbuktu-host.1417: P 17:32(15) ack 1 win 17520 (DF)

22:47:15.675746 Timbuktu-host.1417 > timb-client.1543: . ack 32 win 17489 (DF)

n  
n  
n

22:47:19.871644 Timbuktu-host.1417 > timb-client.1543: P 17471:17478(7) ack 88 win 17433 (DF)

22:47:19.965540 timb-client.1543 > Timbuktu-host.1417: R 3257468433:3257468433(0) win 0 (DF) **[The client is done with the session, tears down session]**








Only after the negotiations are complete, is TCP used. -4 (Netopia FAQ Article)


One of the interesting aspects of the session startup is that it relies on the User Datagram Protocol (UDP) for the exchange of user credentials prior to the establishment of the communications channel. Using UDP for exchange of user credentials is rather counter intuitive:

- Sender and recipient do not keep any information about the state of the communication session between the two hosts
  - UDP Simply provides best-efforts delivery; No guarantee that data is delivered reliably or in order
  - Endpoints do not maintain state information about the communication
  - UDP data is sent and received on a packet-by-packet basis
  - Datagrams must not be too big, because if they must be fragmented, some pieces might get lost in transit
  - In addition, the use of UDP raises **additional security concerns**:
    - When a socket receives data on a UDP port, it will receive packets sent to it by any host, whether it is participating in the application or not
    - This possibility can present a security problem for some applications that do not distinguish between expected and unexpected packets (read: buffer overflow)
    - For these reasons, many network firewall administrators block UDP data from being sent to a protected host from outside the security perimeter
- 5(Australian National University)  
-6(Levier, Laurent)

The exchange of credentials is not particularly time sensitive, nor bandwidth intensive, so the selection of UDP to accomplish this communication is a mystery.


## Issue: Stealth Observation or Control of Timbuktu Host Workstations.

Normal (Not connected)	Timbuktu Connection Type	Active Session	Greyed-out (inactive)
	<b>Control</b>		
	<b>Observe</b>		
	<b>File Exchange</b>		

A Timbuktu Host can be observed or controlled with very little notification (to the individual who is using the Host machine.) There is a small Timbuktu status icon located in the lower right corner of the System Tray. This icon is used to display the current status of the Timbuktu software. When the Timbuktu Host is not being controlled, Observed, Exchanging files, etc, the Timbuktu icon appears to be a small computer. 

When the Timbuktu Host is **Observed** by a remote client machine, the Normal icon is alternated with the Observe icon. The Timbuktu Status flips between the Normal icon and Observe icon about every five to seven seconds.

When the Timbuktu Host is **Controlled** by a remote client machine, the Normal icon is alternated with the Control icon.

After a Timbuktu session is complete, the Timbuktu status icon alternates between the Normal icon and a "greyed-out" version of the icon based on the function(s) used during the remote session. For example: 

Because Timbuktu can not be configured to splash a large notification message/warning across the screen in a remote control / observe / file exchange situation, there is a risk of an attacker or rogue administrator controlling or observing workstations without the end-user being aware of it.

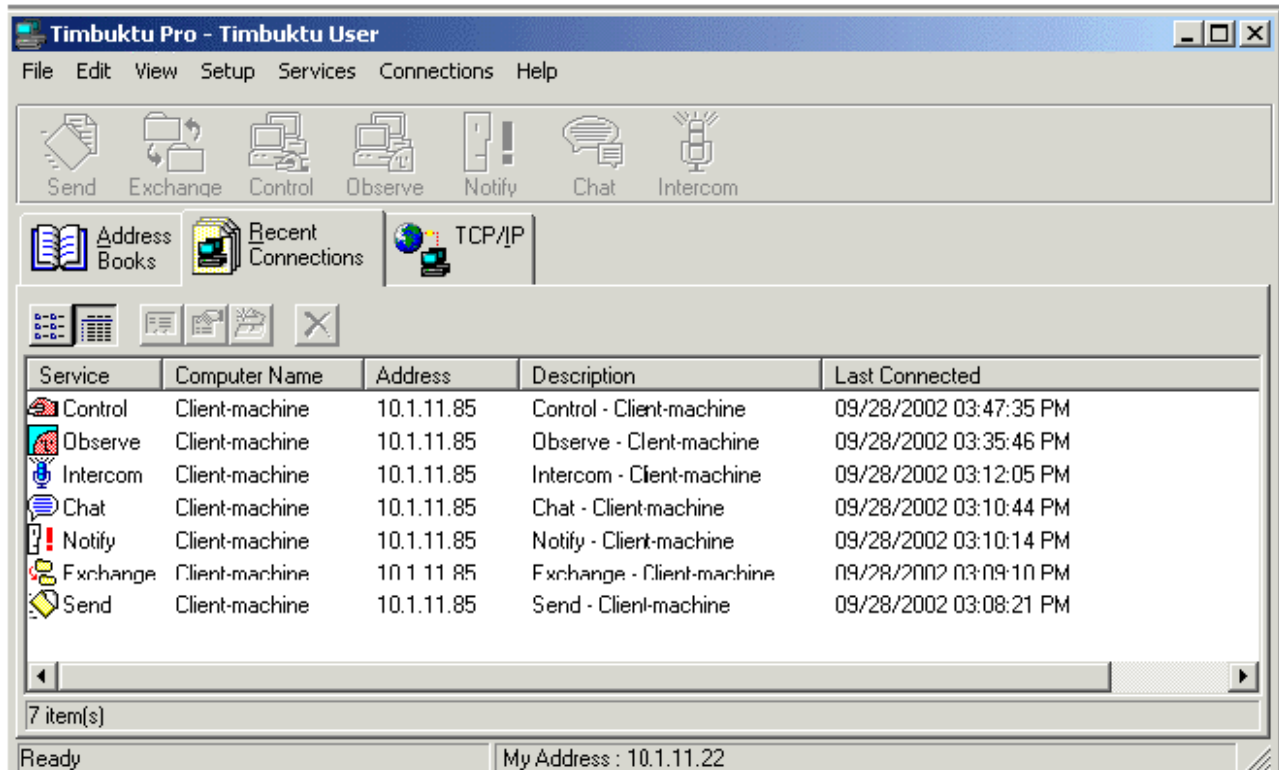
### Potential Attack:

Many users may never notice or understand the significance of a tiny little icon on their task bar changing. Some end users hide the task bar altogether in order to obtain the maximum computer desktop real estate.

Consider the potential damage to a corporation if proprietary business plans or even payroll schedules were being accessed by an individual who was being monitored without their knowledge.

## Logging of Timbuktu Events:

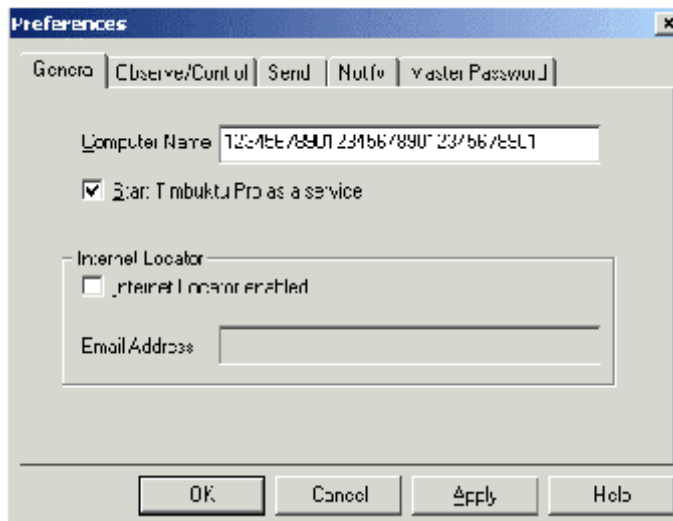
There are some tools available to the Timbuktu host machine to identify remote access. When the user of a Timbuktu Host machine double clicks on the Timbuktu status icon, and selects the Recent Activities Tab, they will see *some* entries in the C:\Program Files\Timbuktu Pro\activity.log file. The GUI presentation is below:



There are some deficiencies with this log:

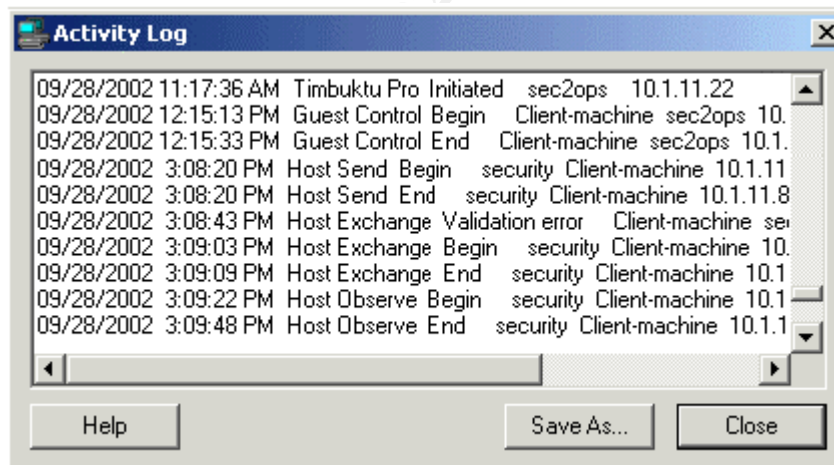
- It only reports activity where the machine functioned as a Timbuktu host
- The Timbuktu Host end-user can delete entries with ease
- An end-user does not require any form of password or authentication to view or delete entries from this screen
- Once a connection entry has been deleted, it can't be recovered (using Timbuktu)
- It would be reasonable to assume that an attacker would delete entries from this file to cover their tracks
- Another point of interest is the data used to populate the field "Computer Name." This name is not necessarily the NetBIOS name of the remote Workstation. The Administrator of the remote workstation can name the Timbuktu Remote Client anything desired.

The maximum Timbuktu Computer name is 31 characters. The following dialog shows where this name can be defined under Setup → Preferences:



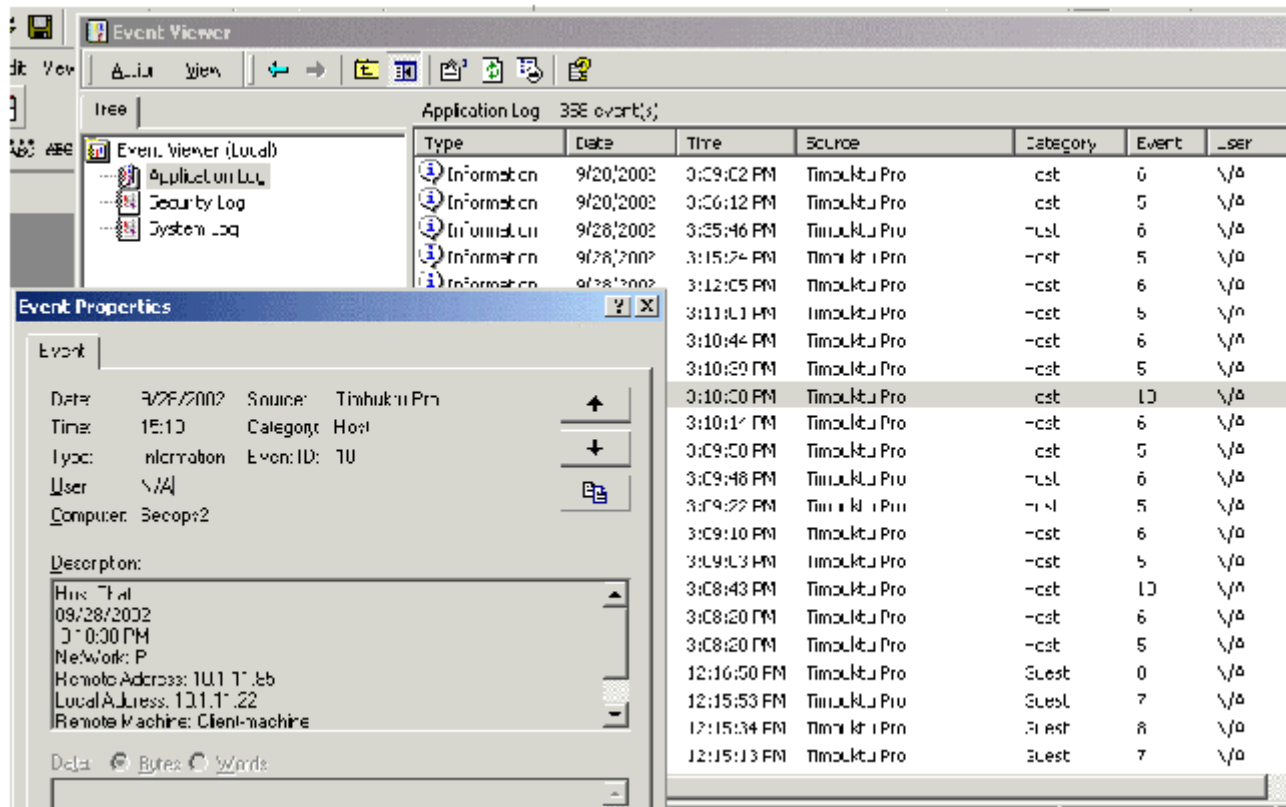
An attacker using Timbuktu to remotely control another workstation could modify their Timbuktu station name in an attempt to mask the origin of unauthorized access.

There are two other places to look for Timbuktu activity as either a Remote Client or a Host: When the user of a Timbuktu machine double clicks on the Timbuktu status icon, and selects Connections → Activity log, the following dialog is displayed:



This is a much more complete log of Timbuktu related events. This dialog will allow the end user to see the entire contents of the C:\Program Files\Timbuktu Pro\activity.log file. By default, the activity.log file is read-only. In a typical installation, only an Administrator or Power-user would be able to delete or manually edit this file. It is reasonable to believe that an attacker would modify or delete this file to hide their activities.

The final tool to review Timbuktu related activities is with the Application Event Viewer:



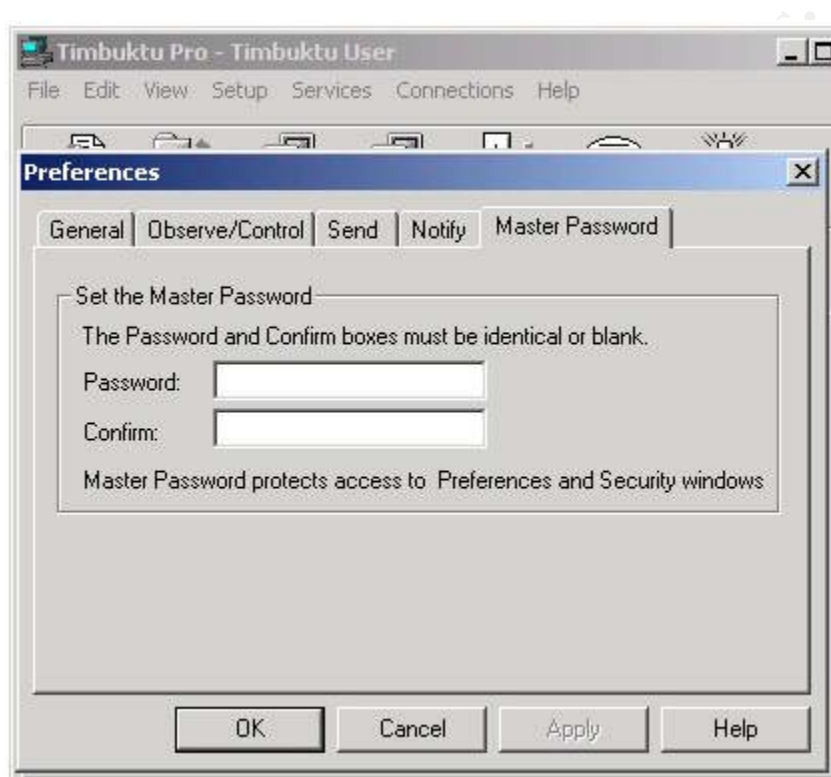
Using the Windows Application Event Viewer, one can obtain very specific and detailed information about Timbuktu related events. Again, it is reasonable to believe that an attacker would clear the event log to cover their tracks.

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## Configuring Access:

Access to the Security Setup screen is controlled by the “Master Password.” If you know the “Master Password” you can add or delete Timbuktu users, or modify their access privileges. –7 (DR.Timbuktu.Database.Insecurity Posting)



This password is stored in an encoded form in the following directory\file:  
C:\Program Files\Timbuktu Pro\tb2.plu.

The password is limited to 15 alpha-numeric characters, and can include spaces and special characters. The password can also include [ALT+Numeric Keypad Sequences] such as [ALT+3333] ♣ .

There are a number of vulnerabilities associated with the tb2.plu file. By default, this file can be read by all users of the computer system. Administrators, Power Users and SYSTEM can write to this file. If Timbuktu is installed on a disk without NTFS or on a Windows 95, 98 or ME workstation, there is no opportunity to create access control lists.

### Issue: Passwords are not Strongly Encrypted

“The password hashes that are generated are not salted. As a result, it is possible to build a dictionary with which the Master Password can be attacked. “

– 8 (Wilson, Rich Security Focus Posting)

Salting password hashes brings randomness, that encoding alone cannot.

- The attacker may pre-compute (offline) encrypted versions of your dictionary. Even if the process is slow, or takes significant storage it may not matter to the attacker, they can just burn a CD or DVD with a database mapping encrypted text to clear-text.
- It's this pre-computation attack which a salt thwarts. A salt makes it impractical to build up a dictionary of encrypted → clear text mappings, because a given clear text has millions of encrypted equivalents. -9(Skoll, David postgres.org posting)

A related problem with the hash for the Master Password, is that it employs the identical algorithm used for encoding passwords for Registered Users. This would allow an attacker to systematically build a dictionary by entering known passwords and recording the encoded output.

The following is an example of the encoding output:

Encoded Password Hexadecimal Value	Plain Text
BCCD E7B6 708A 99D9 6CC5 C49C E31A 1B	aaaaaaaaaaaaaaaa
E612 8A75 6E94 11AA AE95 BEB4 3C33 96	bbbbbbbbbbbbbbbb
6CB0 EDA0 CAD1 5F60 A5B3 CF5C 1284 CD	cccccccccccccccc
B53A 306F 4E58 78F6 2899 7F97 26B0 A8	dddddddddddddddd

### Issue: Registered User Names are Stored in Clear Text

Putting Registered User Names in clear text provides an attacker significant assistance in compromising a system. With knowledge of a valid User name, an attacker can attempt to grind or guess at valid passwords. Historically, passwords have proven to be an extremely weak form of protection from unauthorized access.

-10(Krishna, Arvind Five Steps for Keeping Hackers at Bay)

### Issue: Deletion of C:\Program Files\Timbuktu Pro\tb2.plu Escalates Privileges

One method to defeat the requirement to enter the Master Password is to use the following procedure:

1. Stop the Timbuktu User interface (Tb2pro.exe)
2. Delete C:\Program Files\Timbuktu Pro\tb2.plu
3. Restart the Timbuktu User interface (Tb2pro.exe)
4. Within the Timbuktu User interface, Select Setup, then Preferences
5. Enter a new Master Password of your choice
6. Enter new "Registered Users", "NT Users", or modify Guest User privileges as desired

This procedure will also have the effect of removing or damaging valid Timbuktu Registered User Ids and passwords, which could lead to detection of the attack.

## THE ATTACK:

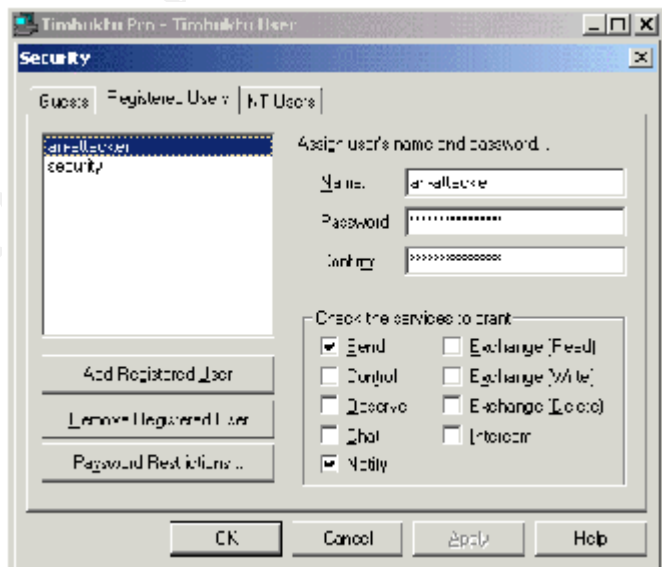
An attacker now has a customized version of the tb2.plu file. From here, the attacker could use a Word Macro virus, Mail Macro Virus, other automated batch or command scripts, or good old fashioned, social engineering to induce victims to place the modified tb2.plu file onto their machines. After the file is placed onto a victim machine, an attacker can take control or “Observe” the machine the next time that Tb2pro.exe starts.

### Issue: Modification of C:\Program Files\Timbuktu Pro\tb2.plu Escalates Privileges Part 1

Let’s consider a scenario where the attacker has a valid Registered User name and a password. But the attacker is not happy because they can only send messages or provide notification to Host computers, and they wish to control Host computers. Modification to the tb2.plu file is a trivial exercise. Armed with nothing more than a hexadecimal editor, an attacker can change 2 bytes, and achieve total power.

Consider the following user setup:

We have a “Registered User” defined with the name of an-attacker. This user does have a password, but has few rights. Only Send and Notify services have been granted.



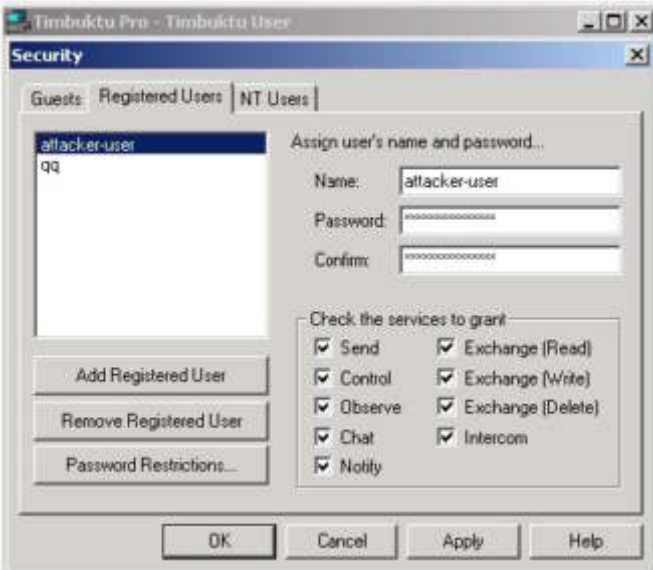
This tb2.plu file excerpt shows the attacker-user with these user rights

Offset	Hexadecimal	ASCII
0000200	0000 0000 0000 0000 0000 0000 FF23 0261	.....#.a
0000210	6E2D 6174 7461 636B 6572 0000 0000 0000	n-attacker.....
0000220	0000 0000 0000 0000 0000 0000 0000 0000	.....
0000230	000F 6ED6 54FB EDCD 95F2 81C2 B221 145E	..n.T.....!^
0000240	1B0F 6ED6 54FB EDCD 95F2 81C2 B221 145E	..n.T.....!^
0000250	1BCA 1C95 3D0F 90BB F8EB F256 CCA5 AF7E	...=......V...~
0000260	6FE3 1661 C900 0000 0000 0000 0000 0000	o..a.....
0000270	0000 0000 0000 0000 0000 0000 0000 0000	.....
0000280	0000 0000 0009 20	.....



Now, let's look at the user setup:

We have now given the attacker-user **all** access rights.



This tb2.plu file excerpt shows the attacker-user with all user rights

```

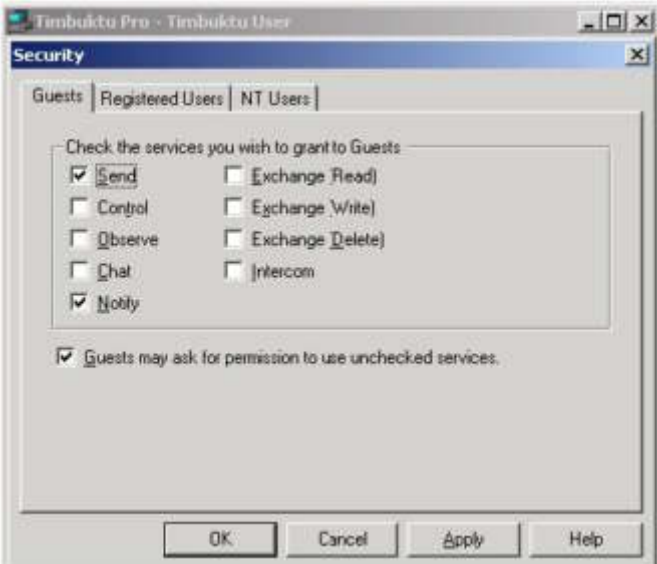
Byte
Offset |-----Hexadecimal-----| |----ASCII-----|
00000200 0000 0000 0000 0000 0000 0000 0000 0261 .....a
00000210 7474 6163 6B65 722D 7573 6572 0000 0000 ttacker-user...
00000220 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000230 000F 224E 35E1 D2DE 3614 7EAC B416 C763 .."N5...6.~....c
00000240 E20F 224E 35E1 D2DE 3614 7EAC B416 C763 .."N5...6.~....c
00000250 E20F DC94 3D00 0000 0000 0000 0000 0000 .....=.
00000260 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000270 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000280 0000 0000 00FF 23 .....#
  
```

**Two bytes in the configuration file was all it took to take a Registered User from few privileges to all privileges.** Due to the nature of the file, an attacker can modify the hexadecimal contents 118 bytes from the start of the clear text user name, and gain all privileges.

**Issue: Modification of C:\Program Files\Timbuktu Pro\tb2.plu Escalates Privileges Part 2**

Earlier, the reader was cautioned about privileges granted to Guests. Using concepts from the previous issue, an attacker could modify the tb2.plu file to escalate privileges of Guests (and in effect, **everybody**.)

In this example, Guests have been given the ability to Send messages, notify the host, and request additional privileges. As a result, **all** users have at least these privileges.

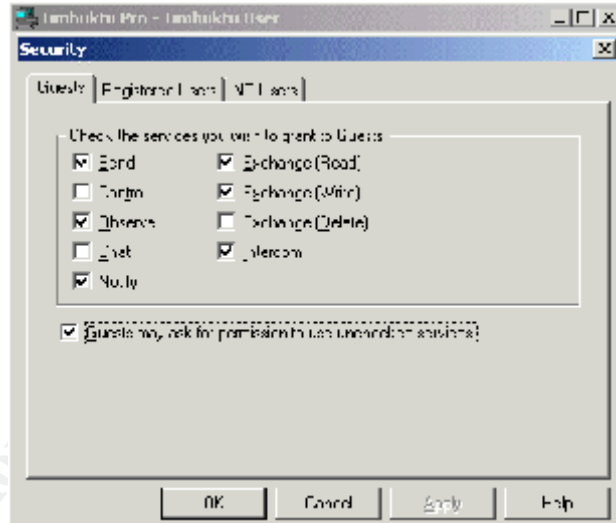


This tb2.plu file excerpt shows the Guest Account with minimal user rights

**Byte**

```
Offset |-----Hexadecimal-----| |----ASCII-----|
000000A0 DB09 6001 3C54 656D 706F 7261 7279 2047 ..`.<Temporary G
000000B0 7565 7374 3E00 0000 0000 0000 0000 0000 uest>.....
```

In this example, additional rights have been granted to Guests.



This tb2.plu file excerpt shows the Guest Account with the added rights

**Byte**

```
Offset |-----Hexadecimal-----| |----ASCII-----|
000000A0 DB7D 6201 3C54 656D 706F 7261 7279 2047 .}b.<Temporary G
000000B0 7565 7374 3E00 0000 0000 0000 0000 0000 uest>.....
```

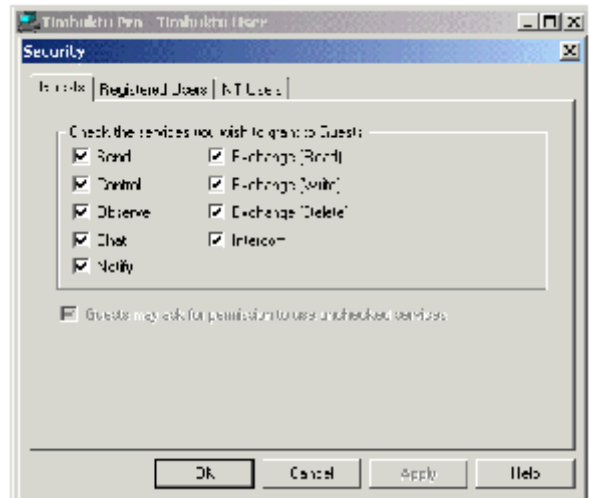
**Two bytes in the configuration file was all it took to modify Guest Privileges.** Look at Hex Offset address 0x000000A1 and 0x000000A2. What happens if the attacker modifies the data at these addresses in a manner similar to the previous example?

**Byte**

```
Offset |-----Hexadecimal-----| |----ASCII-----|
000000A0 DBFF 2301 3C54 656D 706F 7261 7279 2047 ..#.<Temporary G
000000B0 7565 7374 3E00 0000 0000 0000 0000 0000 uest>.....
```

By simply placing the values 0xFF, 0x23 beginning at Hex Offset address 0x000000A1 an attacker can grant the following privileges to everybody.

Results of the modification as seen in the Timbuktu graphical user interface.



## Issue: Modification of C:\Program Files\Timbuktu Pro\tb2.plu Escalates Privileges Part 3

What if the attacker wanted to be more subtle? Deleting the tb2.plu file with the associated Registered Users and privileges might be detected when an authorized person attempted to access a Timbuktu host that had been compromised. The following is an example of a tb2.plu file with two users defined.

```

Byte
Offset  |-----Hexadecimal-----| |----ASCII-----|
00000000 0300 0300 0F67 8651 4C0D 4BB8 DBF9 1253 .....g.QL.K....S
00000010 9398 882B 0000 0000 0100 0000 0000 0000 ...+.....
00000020 0000 0000 0000 0000 0500 003C 4775 6573 .....<Gues
00000030 743E 0000 0000 0000 0000 0000 0000 0000 t>.....
00000040 0000 0000 0000 0000 0000 0000 000F A241 .....A
00000050 0543 9663 E032 9065 46F6 9F6C DB0F A241 .C.c.2.eF..l...A
00000060 0543 9663 E032 9065 46F6 9F6C DB00 0000 .C.c.2.eF..l....
00000070 000F A241 0543 9663 E032 9065 46F6 9F6C ...A.C.c.2.eF..l
00000080 DB0F A241 0543 9663 E032 9065 46F6 9F6C ...A.C.c.2.eF..l
00000090 DB0F A241 0543 9663 E032 9065 46F6 9F6C ...A.C.c.2.eF..l
000000A0 DB01 4001 3C54 656D 706F 7261 7279 2047 ..@.<Temporary G
000000B0 7565 7374 3E00 0000 0000 0000 0000 0000 uest>.....
000000C0 0000 0000 0000 0000 0000 0000 0000 0000 .....
000000D0 0000 0000 0000 0000 0000 0000 0000 0000 .....
000000E0 0000 0000 0000 0000 0000 0000 0000 0000 .....
000000F0 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000100 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000110 0000 0000 0000 0000 0000 0F60 063C 4174 .....`.<At
00000120 7465 6E64 6564 2041 6363 6573 733E 0000 tended Access>..
00000130 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000140 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000150 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000160 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000170 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000180 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000190 0000 00FE 0302 616E 2D61 7474 6163 6B65 .....an-attacke
000001A0 7200 0000 0000 0000 0000 0000 0000 0000 r.....
000001B0 0000 0000 0000 0000 0F90 BBF8 EBF2 56CC .....V.
000001C0 A5AF 7E6F E316 61C9 0F90 BBF8 EBF2 56CC ..~o..a.....V.
000001D0 A5AF 7E6F E316 61C9 1C14 953D 0000 0000 ..~o..a....=....
000001E0 0000 0000 0000 0000 0000 0000 0000 0000 .....
000001F0 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000200 0000 0000 0000 0000 0000 0000 0000 0273 .....s
00000210 6563 7572 6974 7900 0000 0000 0000 0000 ecurity.....
00000220 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000230 000F 82B6 0D61 EA5E B6F4 F634 1C0E 3F43 .....a.^...4..?C
00000240 5A0F 82B6 0D61 EA5E B6F4 F634 1C0E 3F43 Z....a.^...4..?C
00000250 5A1C 1495 3D0F AED0 2234 9B8B 3075 626F Z...=..."4..0ubo
00000260 369F ACC6 D600 0000 0000 0000 0000 0000 6.....
00000270 0000 0000 0000 0000 0000 0000 0000 0000 .....
00000280 0000 0000 00FF 23 .....#

```

The Registered User names can be seen beginning at offset 0x0196 “an-attacker”, and at offset 0x020F “security.” An attacker can modify their own password, then examine the file again, looking for changes. Close observation of changes to the file when the password changes show that one can always predict where the current password will be located in

relation to the Registered User name. The current password always starts 35 bytes after the beginning of the Registered User name. In addition, the encoded password is always 15 bytes in length, even if the password itself is only one character long. Further experimentation shows that the Master Password field always begins at offset 0x05. Therefore, an attacker could perform the following steps to gain Master Password level access:

1. Use the Timbuktu application to generate an encoded password for a known password
2. Make a note of the original encoded Master Password beginning at offset 0x05 for a length of 15 bytes
3. Stop the Timbuktu User interface (Tb2pro.exe)
4. Modify C:\Program Files\Timbuktu Pro\tb2.plu, placing the known encoded password at offset 0x05
5. Start Timbuktu User interface (Tb2pro.exe)
6. Enter the known password
7. Enter new "Registered Users", "NT Users", or modify Guest User privileges as desired
8. Stop the Timbuktu User interface (Tb2pro.exe)
9. Modify C:\Program Files\Timbuktu Pro\tb2.plu, placing the original encoded password (of the valid Master Password) at offset 0x05.

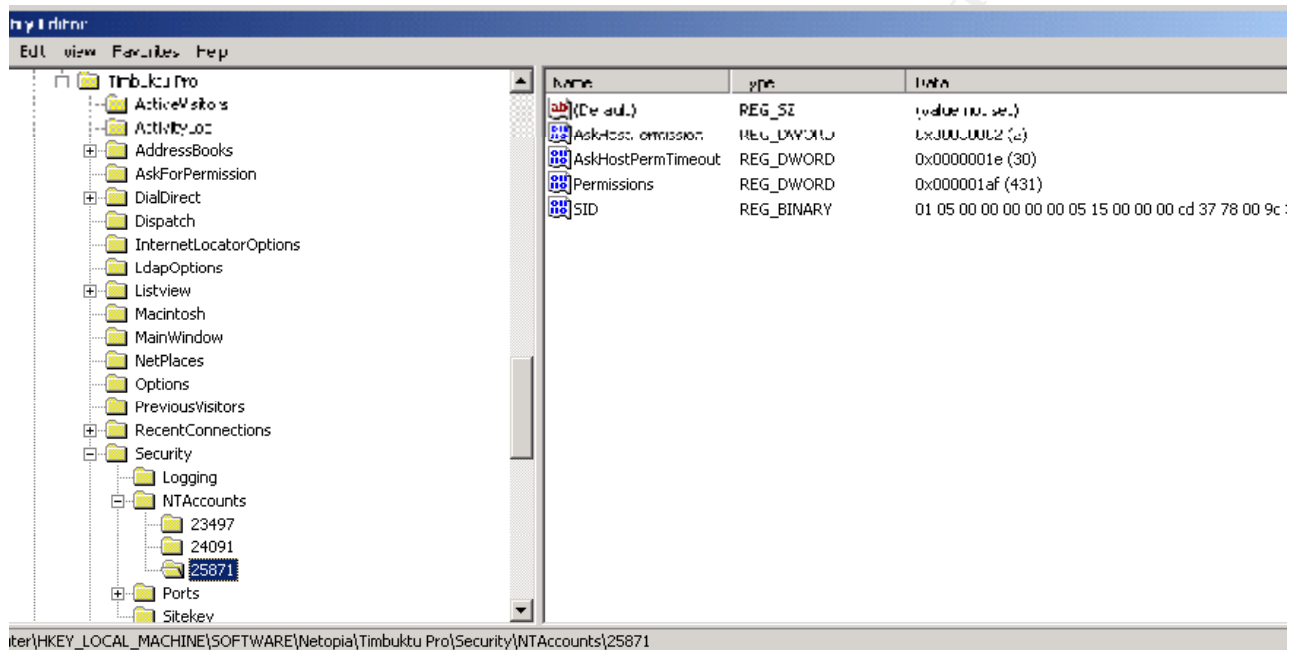
#### **Issue: Modification of C:\Program Files\Timbuktu Pro\tb2.plu Escalates Privileges Part 4**

What if the attacker wanted to be subtle but did not want to have to mess around with copying an encoded password to a known password? There is an even easier way of defeating the "security" of the Master Password. Looking at the example tb2.plu file, one notices several patterns. One of the patterns is that it appears that the encoded password is repeated multiple times for each Registered User Id. It appears that this repeated hash is the mechanism that Timbuktu uses to detect if a user attempts to reuse a password within three generations. The other pattern that can be seen is a sequence of 0x0F preceding each encoded password. This sequence also precedes the encoded Master Password. Testing has shown that if the sequence of 0x0F at offset 0x04 is modified to be 0x00, one is no longer prompted to enter the Master Password in order to add or modify Timbuktu users.

1. Stop the Timbuktu User interface (Tb2pro.exe)
2. Modify C:\Program Files\Timbuktu Pro\tb2.plu, placing value 0x00 at offset 0x04
3. Start Timbuktu User interface (Tb2pro.exe)
4. Enter new "Registered Users", "NT Users", or modify Guest User privileges as desired
5. Stop the Timbuktu User interface (Tb2pro.exe)
6. Modify C:\Program Files\Timbuktu Pro\tb2.plu, placing value 0x0F at offset 0x04

## Issue: Access to NT Registry may allow attacker to Modify or Destroy Timbuktu NT User Account information

Timbuktu NT User(and Group) membership information is stored in the registry. If an attacker has the ability to make manual changes to the registry using a tool such as Regedit, they may modify permissions to escalate privileges. Although this type of attack is possible, it is more difficult to perform successfully, as the attacker would need to understand the mapping of the NT User/Group information to the SID values. A more likely possibility is that the attacker would simply delete keys located below `\HKEY_LOCAL_MACHINE\SOFTWARE\Netopia\Timbuktu Pro\Security\NTAccounts`.



**Issue: Registered User Ids are sent in clear text during initial session establishment.** Although this might not seem to be a very large vulnerability, it can lead to a severe breach in security. Organizations or individuals may use Timbuktu software on Windows devices that are directly connected to the Internet.

Without launching into a lengthy discussion on the perils of leaving unprotected Windows machines connected to the Internet, let's just focus on the Timbuktu access issues.

It was previously identified that all Timbuktu session negotiation and credential exchange occurs on UDP port 407. Therefore it is reasonable to expect that a high percentage of computers that have UDP port 407 open, function as Timbuktu hosts. If an attacker can get valid user credentials for a particular Timbuktu host, at least half of the work of compromising a system has been completed.

The following is a Ethereal packet capture of a Timbuktu session negotiation:

No.	Time .	Source	Destination	Protocol	Info
11	20.339328	timb-client	Timbuktu-host	UDP	Source port: 1549 Destination port: 407
12	20.344371	Timbuktu-host	timb-client	UDP	Source port: 407 Destination port: 1549
13	26.242745	timb-client	Timbuktu-host	UDP	Source port: 1550 Destination port: 407
14	26.245362	Timbuktu-host	timb-client	UDP	Source port: 407 Destination port: 1550
15	26.243160	timb-client	Timbuktu-host	UDP	Source port: 1551 Destination port: 407
16	26.253665	Timbuktu-host	timb-client	UDP	Source port: 407 Destination port: 1551
17	26.258709	timb-client	Timbuktu-host	TCP	1552 > 1417 [SYN, Seq=2728227574, Win=0, Len=0] window=1638
18	26.258525	Timbuktu-host	timb-client	TCP	1417 > 1552 [SYN, ACK] Seq=2724219529, Ack=2728227574, Win=0, Len=0
19	26.258572	timb-client	Timbuktu-host	TCP	1552 > 1417 [ACK] Seq=2728227575, Ack=2724219530, Win=0, Len=0
20	26.253943	timb-client	Timbuktu-host	TCP	1552 > 1417 [PSH, ACK] Seq=2728227575, Ack=2724219530, Win=0, Len=0

In frames 11 through 16, the client is talking to destination port UDP 407. In frames 17 through 19 the three-way TCP handshake is completed. The actual remote control session begins in frame 20.

Let's take a look at the traffic exchanged in frames 11 through 16 and see if there is any information present that would be interesting to an attacker.

In frame 11, there is no obviously useful information

No.	Time .	Source	Destination	Protocol	Info
11	20.339328	timb-client	Timbuktu-host	UDP	Source port: 1549 Destination port: 407
12	20.344371	Timbuktu-host	timb-client	UDP	Source port: 407 Destination port: 1549

0000	00 80 c7 a8 e8 75 00 d0	59 39 d9 be 08 00 45 00	.....u..Y9....E.
0010	00 6e b6 d0 00 00 80 11	59 07 0a 01 0b 55 0a 01	.n.....Y....U..
0020	0b 51 06 0d 01 97 00 5a	84 06 00 25 00 22 00 01	.Q.....Z ...%. "...
0030	00 0f 03 08 00 05 00 00	06 00 00 00 00 00 00 00	.....
0040	00 00 00 00 00 00 14 ec	06 00 0f 6b 1e 00 98 04	.....k....
0050	00 00 17 7d db 77 98 04	00 00 00 00 00 00 28 ec	...}.w.. ....(.
0060	06 00 dd 7c db 77 3c f7	06 00 00 00 00 00 04 20	... .w<. ....
0070	40 00 14 ec 06 00 80 52	21 03 b8 f3	@.....R !....

Let's look at frame 12.

No.	Time .	Source	Destination	Protocol	Info
11	20.339320	timb-client	timbuktu-host	UDP	Source port
12	20.344371	Timbuktu-host	timb-client	UDP	Source port

0000	00 d0 59 39 d9 be 00 80	c7 a8 e8 75 08 00 45 00	..Y9....	...u..E.
0010	00 aa 0b a4 00 00 80 11	03 f8 0a 01 0b 51 0a 01	.....	.....Q..
0020	0b 55 01 97 06 0d 00 96	f6 a9 00 25 00 00 00 00	.U.....	...%....
0030	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0040	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0050	0f 53 45 43 55 52 49 54	59 2d 32 4b 2d 54 45 53	.SECURIT	Y-2K-TES
0060	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0070	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0080	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0090	00 00 01 03 90 0f bb 55	00 41 94 82 00 00 00 00	.....U	.A.....
00a0	00 08 00 0f 01 00 01 00	01 00 00 03 00 00 00 05	.....	.....
00b0	00 00 00 00 00 00 00 00		.....	.....

Now we have something of interest. Take a look at Hex offset 0x0051. That is the Timbuktu name of the Host machine (SECURITY-2K-TES). An attacker will use this and any other available information to develop a victim vulnerability profile.

Frame 13: In this frame, there is no obviously useful information.

No.	Time .	Source	Destination	Protocol	Info
12	20.344371	Timbuktu-host	timb-client	UDP	Source port
13	26.240745	timb-client	Timbuktu-host	UDP	Source port

0000	00 80 c7 a8 e8 75 00 d0	59 39 d9 be 08 00 45 00	.....u..	Y9....E.
0010	00 6e b6 ed 00 00 80 11	58 ea 0a 01 0b 55 0a 01	.n.....	X....U..
0020	0b 51 06 0e 01 97 00 5a	d3 6d 00 25 00 22 00 01	.Q.....Z	,m.%."..
0030	00 0f 03 08 00 05 00 00	22 00 4e 02 25 00 82 00	.....	".N.%...
0040	00 00 00 00 00 00 00 00	07 00 00 31 bc 01 00 31	.....	...1...1
0050	bc 01 00 00 00 00 f8 eb	fd 7f 00 00 00 00 00 00	.....	.0.....
0060	00 00 00 00 00 00 5c eb	06 00 98 04 00 00 00 00	.....\	.....
0070	07 00 f8 eb fd 7f 00 00	00 00 70 eb	.....0..	..p.

How about Frame 14?

No.	Time .	Source	Destination	Protocol	Info
13	26.240745	timb-client	Timbuktu-host	UDP	Source port
14	26.246382	Timbuktu-host	timb-client	UDP	Source port

0000	00 d0 59 39 d9 be 00 80	c7 a8 e8 75 08 00 45 00	..Y9....	...u..E.
0010	00 aa 0b a5 00 00 80 11	03 f7 0a 01 0b 51 0a 01	.....	.....Q..
0020	0b 55 01 97 06 0e 00 96	f6 a8 00 25 00 00 00 00	.U.....	...%....
0030	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0040	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0050	0f 53 45 43 55 52 49 54	59 2d 32 4b 2d 54 45 53	.SECURIT	Y-2K-TES
0060	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0070	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0080	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	.....	.....
0090	00 00 01 03 90 0f bb 55	00 41 94 82 00 00 00 00	.....U	.A.....
00a0	00 08 00 0f 01 00 01 00	01 00 00 03 00 00 00 05	.....	.....
00b0	00 00 00 00 00 00 00 00		.....	.....

This seems to be nearly a repeat of frame 12. Other than the repeated host name, there is no obviously useful information

How about Frame 15?

No.	Time .	Source	Destination	Protocol	Info
14	26.248382	Timbuktu-host	timb-client	UDP	Source p
15	26.248460	timb-client	Timbuktu-host	UDP	Source p
*****					
*****					
0000	00 80 c7 a8 e8 75 00 d0	59 39 39 be 08 00 45 00	.....U..	Y9....E.	
0010	00 cc b6 f4 00 00 80 11	58 85 3a 01 0b 55 0a 01	.....X....U..		
0020	0b 51 06 0f 01 97 00 b8	f1 25 30 23 07 22 03 00	.Q.....	.%.#."..	
0030	90 0f bb 55 00 41 94 82	61 af 21 45 8b a7 d1 57	...U.A..	a.!E...W	
0040	02 01 1a 61 62 63 64 65	66 67 68 69 6a 6b 6c 6d	...abcde	fg hijklm	
0050	6e 6f 70 71 72 73 74 75	76 77 78 79 7a 00 00 00	nopqrstu	vwxyz...	
0060	00 00 00 00 00 00 00 00	00 00 30 00 00 00 00 00	.....	.....	
0070	00 00 00 00 00 00 00 00	00 00 30 00 0e 43 6c 69	.....	.....cli	
0080	65 6e 74 2d 6d 61 63 68	69 6e 35 00 00 00 00 00	ent-mach	ine.....	
0090	00 00 00 00 00 00 00 00	00 00 30 00 00 00 00 00	.....	.....	
00a0	00 00 00 00 00 00 00 00	00 00 30 00 00 00 00 00	.....	.....	
00b0	00 00 00 00 00 00 00 00	00 00 30 00 00 00 00 01	.....	.....	
00c0	00 0f 34 4e 30 3b 00 00	00 00 30 00 00 00 00 00	..4N0;..	.....	
00d0	00 00 00 00 00 00 00 00	00 00	.....	..	

Here we can clearly see two pieces of valuable information. At hex offset 0x0043 we can see the Timbuktu user name. The name is unusual in that it is “abcdefghijklmnopqrstuvwxy”, but I wanted to use an obvious character string. Also, at hex offset 0x007D, we can see the name of the Timbuktu client workstation. The name is “Client-machine.” Now the attacker has two more pieces of information.; a valid user name and the name of Timbuktu machine that a valid client may come from.

Here is Frame 16:

No.	Time .	Source	Destination	Protocol	Info
15	26.248460	timb-client	Timbuktu-host	UDP	Source
16	26.256965	Timbuktu-host	timb-client	UDP	Source
*****					
*****					
0000	00 d0 59 39 d9 be 00 80	c7 a8 e8 75 08 00 45 00	..Y9....	...U..E.	
0010	00 48 0b a6 00 00 80 11	04 58 0a 01 0b 51 0a 01	.H.....	.X...Q..	
0020	0b 55 01 97 06 0f 00 34	0a c5 00 23 00 00 00 00	.U.....4	...#....	
0030	01 00 00 0f 05 89 00 00	0b 58 99 3b ff 23 00 00	.....	.X.;.#..	
0040	00 00 00 00 00 00 00 00	00 00 00 00 00 00 04 00	.....	.....	
0050	03 00 10 01 01 00		.....	.....	

No obviously useful information here other than what an acknowledgement of valid credentials looks like.



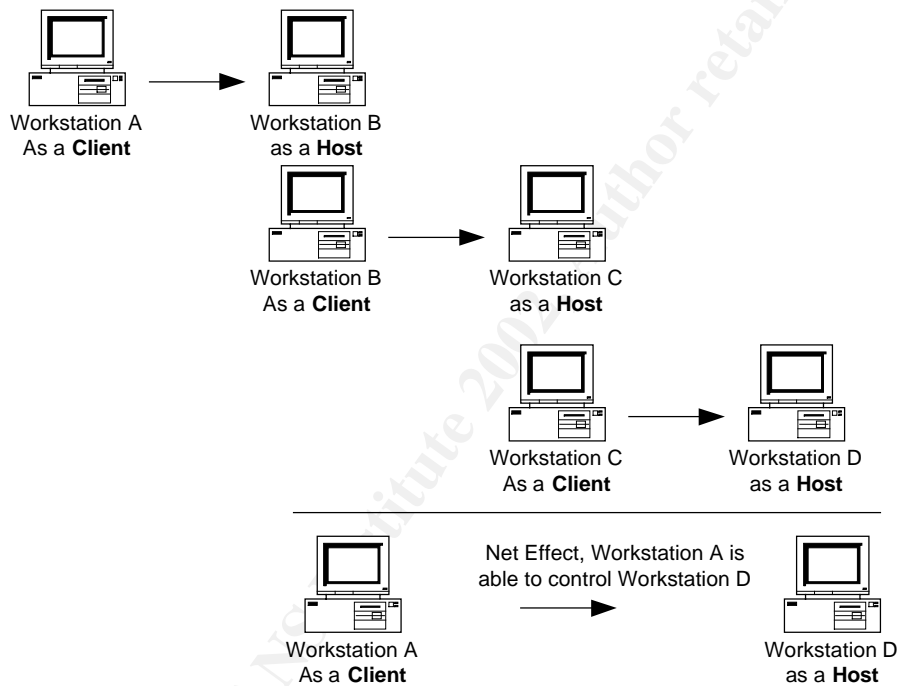
It is important to note that during this exchange, the password was not sent in clear text. Also, the encoded password was not sent in clear text. It is my belief that the client software re-encodes the password prior to sending it from the client to the host.

---

### Security risks Associated with Remote Control Software:

An important feature of Timbuktu is that a device can function as both a Timbuktu host and client simultaneously. For example, Workstation A can connect as a client to Workstation B. Workstation B can then (while under the control of Workstation A) initiate a connection to Workstation C.

While not unique to Timbuktu, the ability to leapfrog functional control from machine to machine can lead to unanticipated security vulnerabilities and exposures.



In the example above, workstation A controls workstation D, even with firewall(s) separating the two machines.

To avoid problems of this type, organizations should follow the principle of least privilege access.

Remote access functionality should not be allowed to span borders of autonomous groups, nor should non-secured devices be allowed to control high secured, highly sensitive servers.

## **Defense in Depth Response to the challenges posed by Timbuktu and Remote access tools in General.**

1. Understand and mitigate the risks
  - Recognize that the use of Remote control software effectively extends the network perimeter to include the Remote clients.
  - Avoid software/technical solutions where one device can simultaneously function as a Remote host and a client.
2. Develop Policy to address the risks
  - Prohibit the use of remote control software implementations that have not been reviewed and approved by the Information Technology Security Department.
  - Educate Help Desk staff and other support personnel of their obligation to obtain consent and inform individuals when their workstations are being accessed or controlled remotely.
  - Prohibit remote control/remote observation software from loading automatically on standard end-user desktops. If help desk or support personnel require remote access to an end-user workstation, require the end-user to start the host software.
  - Educate all employees that any attempt to escalate privileges through the use of remote control software or other means is expressly prohibited, and that individuals found to engage in such activities are subject to termination and potential civil litigation.
  - Develop corporate policy to establish minimum password requirements.
  - Audit remote control passwords to verify that they comply with established corporate policy.
  - Establish policy concerning administrative boundaries or domains for workstations and servers. Prohibit unsecured devices from controlling critical or sensitive devices.
  - Prohibit external hosts from controlling or connecting to internal workstations over the Internet.
  - Develop a Computer Incident Response Team to respond to computer security breaches/incidents.
3. Identify Appropriate Remote Access/Remote Control Technology
  - Use solutions that employ strong encryption for communications.
  - Use solutions that employ strong encryption for locally saved configuration files.
  - Use solutions that employ two-factor authentication.
  - Secure critical configuration files that may be stored on workstations with appropriate access control lists.
4. Monitor The Infrastructure
  - Monitor critical configuration files that may be stored on workstations/servers with Tripwire or other tools to be made aware of unauthorized changes.
  - Use secure, centralized logging for remote access events. Be able to document who was accessing what remote computer when.

- Use network and host based intrusion detection systems to be alerted of abnormal remote access attempts.
  - Perform pro-active network scans on a periodic basis to identify rogue or unauthorized remote control software installations
5. Provide Appropriate Training
- Educate administrators about the risks associated with remote access tools.
  - Provide training to all computer users, but System Administrators in particular about steps to take in the event they believe security has been breached.
  - Train all computer users on need to maintain strong passwords.

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